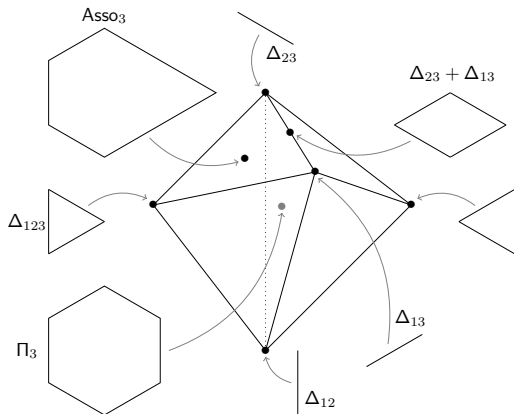


Geometric combinatorics of paths and deformations of convex polytopes

Germain Poullot



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Directeurs / Advisors

Arnau PADROL

Vincent PILAUD

Jury

Fu LIU

Jesús DE LOERA

Martina JUHNKE-KUBITZKE

Lionel POURNIN

Frédéric MEUNIER

Vic REINER

1 What is “Combinatorics of Polytopes”?

2 Generalized permutahedra

- Deformations
- Submodular Cone
- Ongoing work

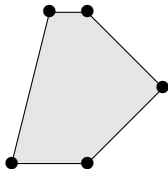
3 Max-slope Pivot Polytopes

- Max-slope pivot rule
- Poset of slopes
- Pivot rule polytope of products of simplices

What is “Combinatorics of Polytopes”?

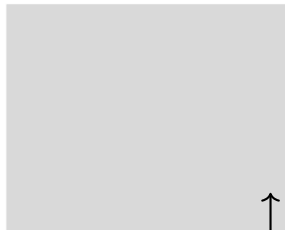
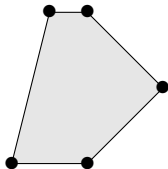
Definition

Polytope: convex hull of finitely many points in $\mathbb{R}^n$



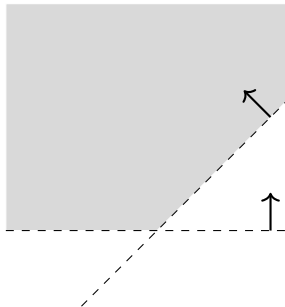
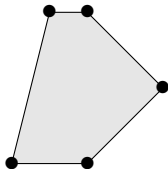
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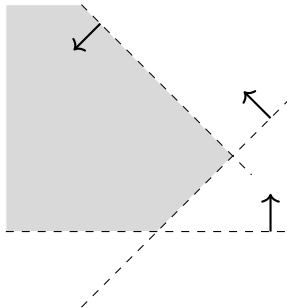
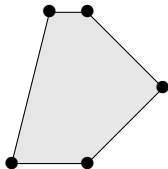
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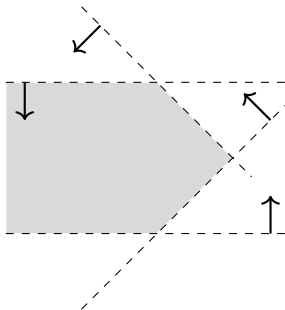
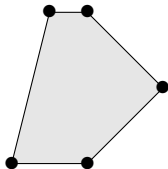
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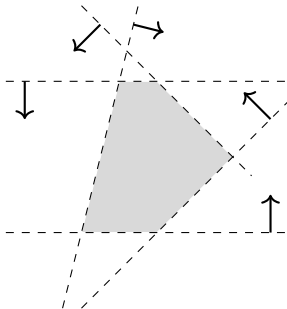
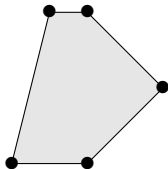
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Representing polytopes



Tetrahedron
Fire



Hexahedron
Earth



Octahedron
Air



Dodecahedron
the Universe



Icosahedron
Water

Representing polytopes



Tetrahedron
Fire



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Representing polytopes



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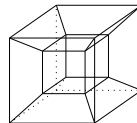
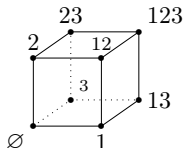
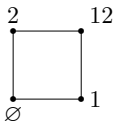
Octahedron
Air



Dodecahedron
the Universe

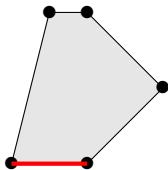


Icosahedron
Water



Definition

Face: $P^c := \{ \mathbf{x} \in \mathbb{R}^n ; \langle \mathbf{x}, \mathbf{c} \rangle = \max_{\mathbf{y} \in P} \langle \mathbf{y}, \mathbf{c} \rangle \}$

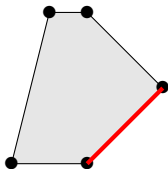


P



Definition

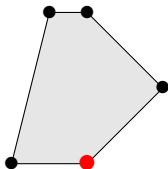
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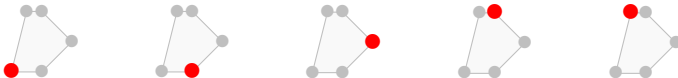


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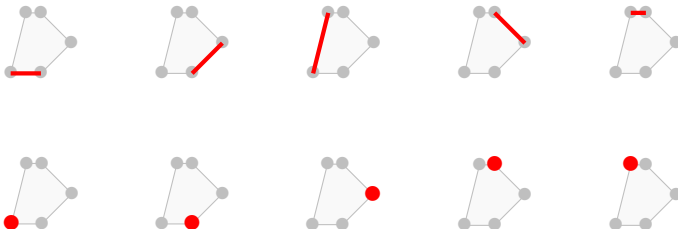
Face lattice

Face lattice: poset of inclusions of faces



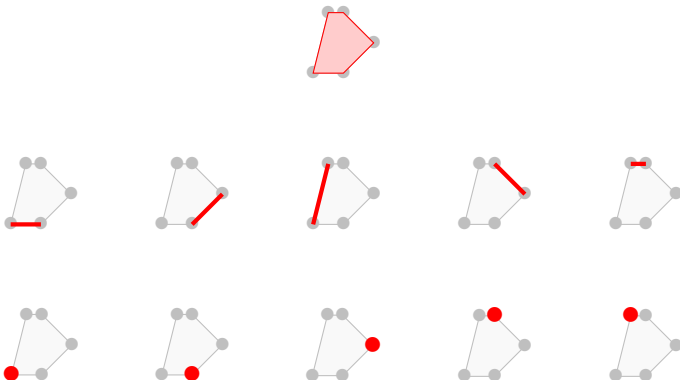
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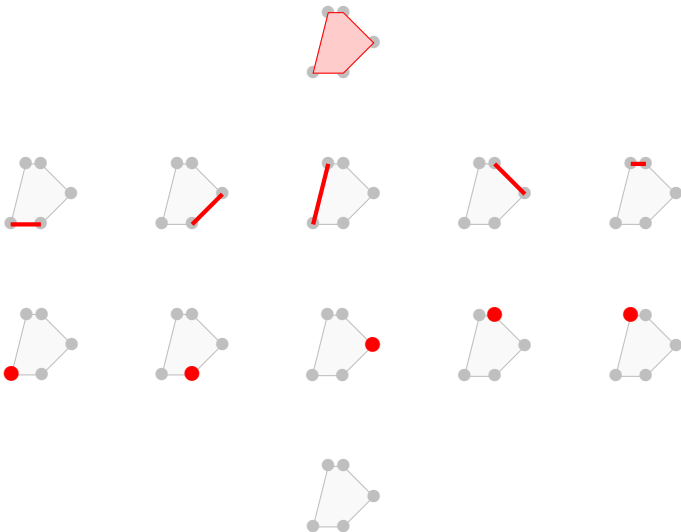
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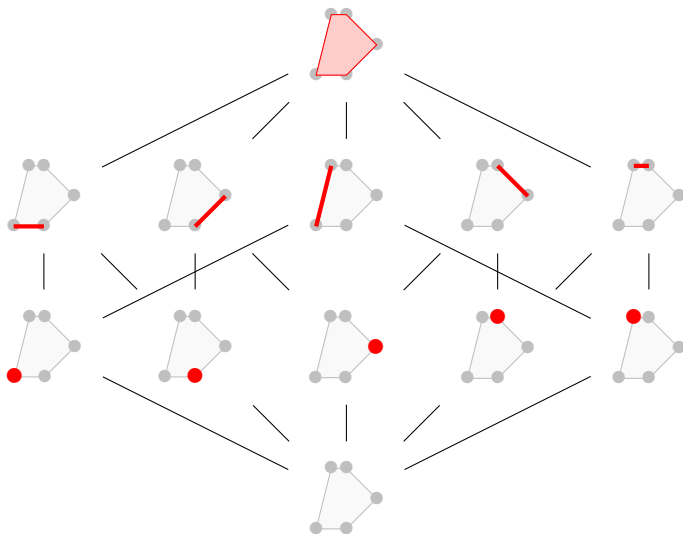
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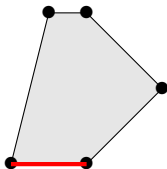
Face lattice: poset of inclusions of faces



Definition

Normal cone of a face F : $\mathcal{N}_P(F) := \{\mathbf{c} ; P^c = F\}$

Normal fan $\mathcal{N}_P$: collection of $\mathcal{N}_P(F)$ for F face of P



P

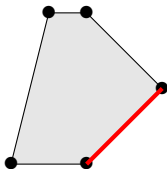


$\mathcal{N}_P$

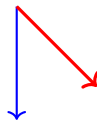
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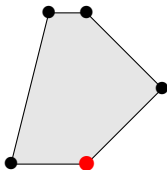


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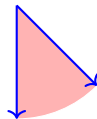
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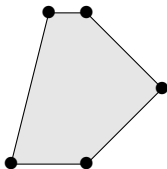


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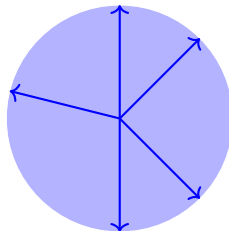
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Combinatorics of Polytopes

One way: Take a polytope $\rightarrow$ combinatorial info (e.g. face lattice)

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Example (Permutahedron)

$$\Pi_n = \text{conv} \left\{ \begin{pmatrix} \sigma(1) \\ \vdots \\ \sigma(n) \end{pmatrix} ; \sigma \text{ permutation of } \{1, \dots, n\} \right\}$$

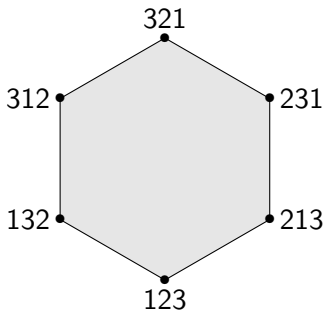
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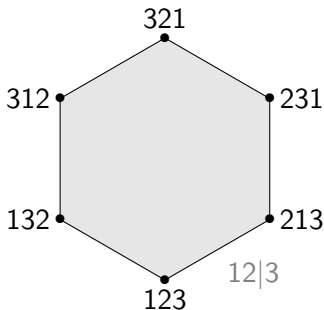
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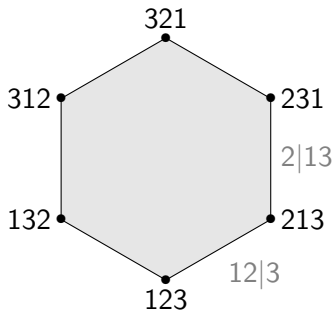
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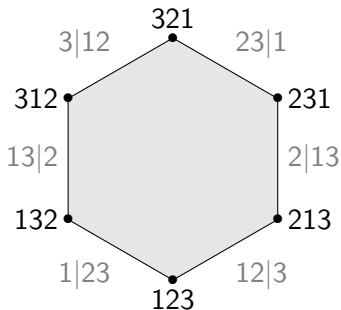
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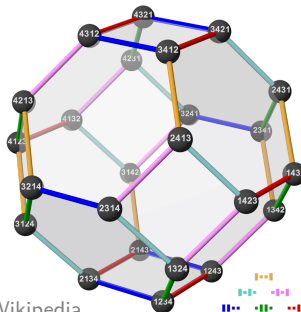
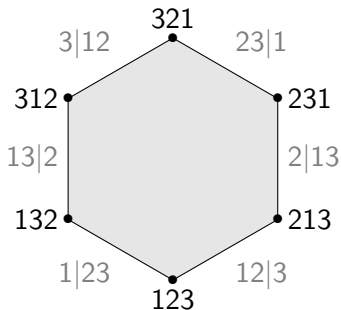
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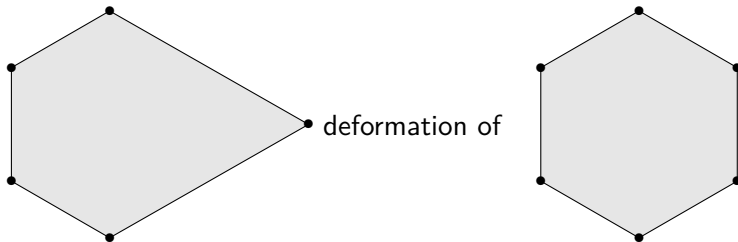
Wikipedia

Generalized permutahedra

Coarsening: Choose maximal cones and merge them

Definition

Q is a *deformation* of P iff $\mathcal{N}_Q$ coarsens $\mathcal{N}_P$.

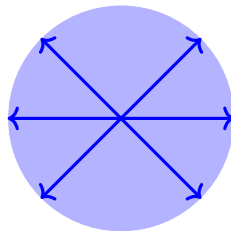
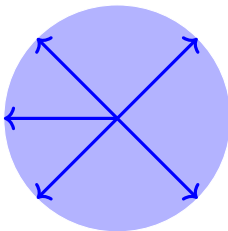
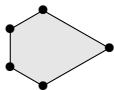


Deformations

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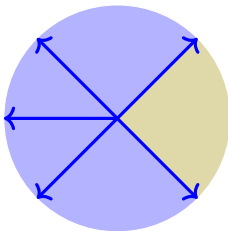
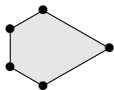


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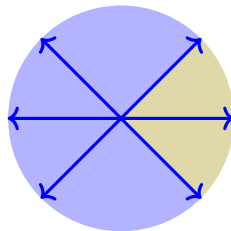
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coarsens



Definition

Braid fan: arrangement of hyperplanes $H_{i,j} := \{\mathbf{x} ; x_i = x_j\}$

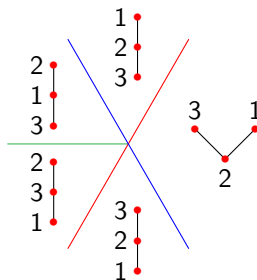
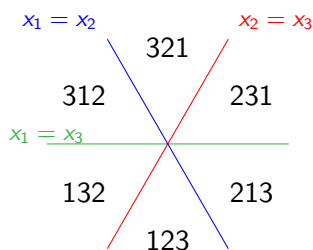
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Generalized permutahedron: deformation of Π_n
i.e. P generalized permutahedron iff $\mathcal{N}_P$ coarsens braid fan



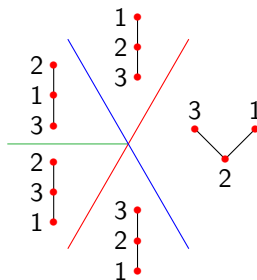
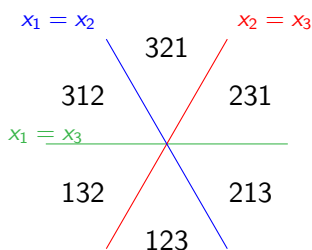
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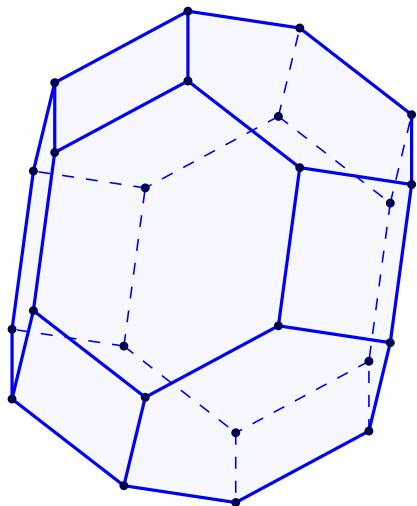
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$\mathcal{P}(P)$: all the posets associated to faces of P

Deformations of Π_4



Permutahedron Π_4

Sequence of deformations of Π_4

Cone of deformations

Minkowski sum: $P + Q = \{\mathbf{p} + \mathbf{q} ; \mathbf{p} \in P, \mathbf{q} \in Q\}$

Theorem

If Q, R deformations of P , then: *for all $\lambda > 0$, λQ deform. of P*
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Definition

Deformation cone: $\mathbb{DC}(P) := \{Q ; Q \text{ deformation of } P\}$ is a cone.

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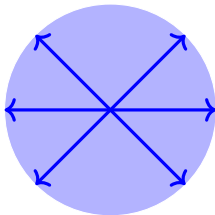
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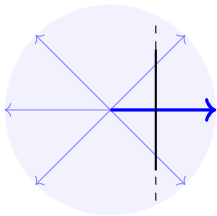
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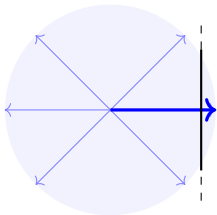
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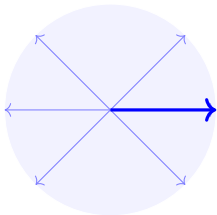
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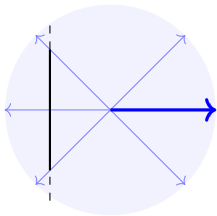
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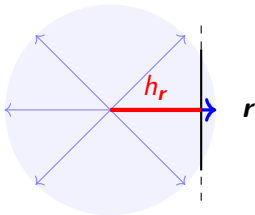
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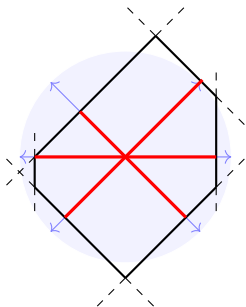
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Parametrization:

height vector:

$$\mathbf{h} = (h_r)_{r \text{ rays}}$$

Cone of deformations

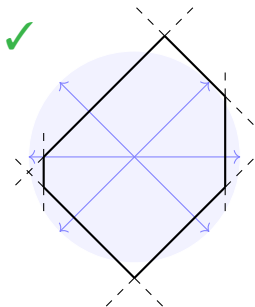
Minkowski sum: $P + Q = \{\mathbf{p} + \mathbf{q} ; \mathbf{p} \in P, \mathbf{q} \in Q\}$

Theorem

If Q, R deformations of P , then: *for all $\lambda > 0$, λQ deform. of P*
 $Q + R$ deform. of P

Definition

Deformation cone: $\mathbb{DC}(P) := \{Q ; Q \text{ deformation of } P\}$ is a cone.



Parametrization:

height vector:

$$\mathbf{h} = (h_r)_{r \text{ rays}}$$

Wall-crossing inequalities:

linear ineq on $\mathbf{h}$

Cone of deformations

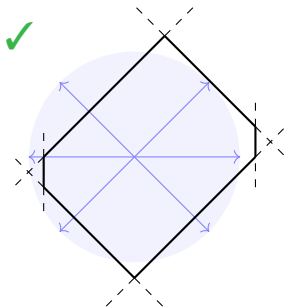
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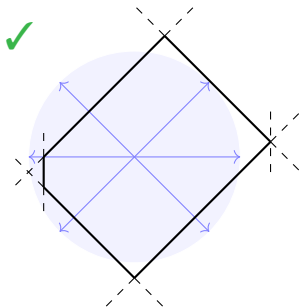
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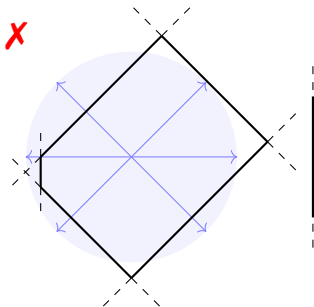
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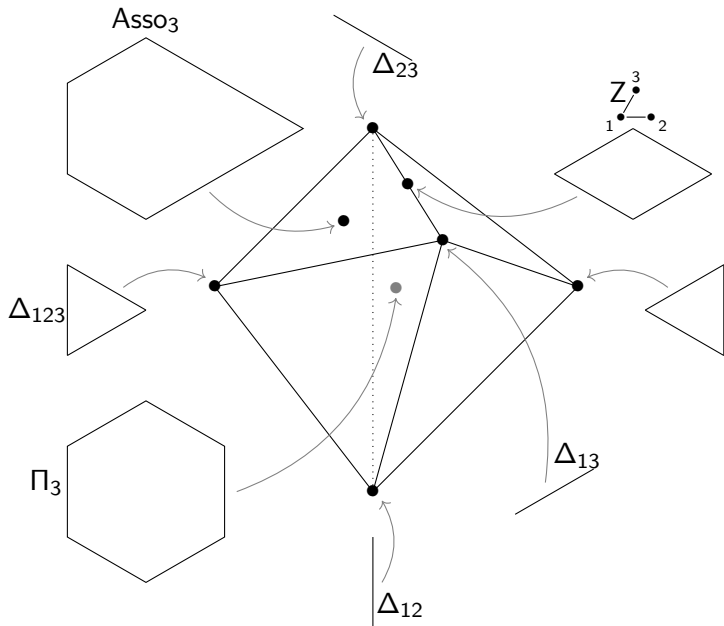
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Definition

Submodular cone: deformation cone of the permutahedron Π_n

	$\mathbb{DC}(\Pi_n)$
Dim (no lineal)	$2^n - n - 1$
# facets	$\binom{n}{2} 2^{n-2}$
# rays	unknown!

Submodular Cone for Π_3



Definition

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Submodular Cone's faces

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If Q deformation of P , then $\mathbb{DC}(Q)$ is a face of $\mathbb{DC}(P)$.

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	$\mathbb{DC}(\Pi_n)$	$\mathbb{DC}(\text{Asso}_n)$
Dim (no lineal)	$2^n - n - 1$	$\binom{n}{2}$
# facets	$\binom{n}{2} 2^{n-2}$	$\binom{n}{2}$
# rays	unknown!	$\binom{n}{2}$
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	$\mathbb{DC}(\Pi_n)$	$\mathbb{DC}(\text{Asso}_n)$	$\mathbb{DC}(Z_G)$	$\mathbb{DC}(N_B)$
Dim (no lineal)	$2^n - n - 1$	$\binom{n}{2}$	N	N
# facets	$\binom{n}{2} 2^{n-2}$	$\binom{n}{2}$	E	E
# rays	unknown!	$\binom{n}{2}$	X	X
		is simplicial!	T	T

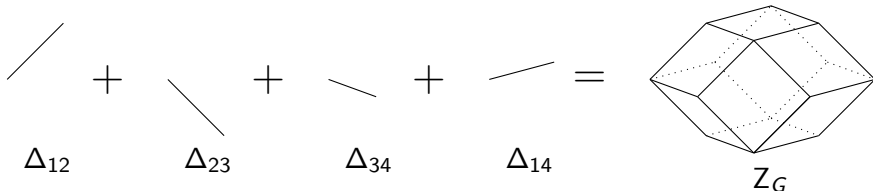
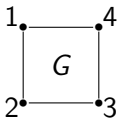
My contribution - Graphical Zonotopes

$G = (V, E)$ a graph, $n = |V|$

Definition

Graphical zonotope $Z_G := \sum_{(i,j) \in E} [\mathbf{e}_i, \mathbf{e}_j]$

Z_G deformation of $\Pi_n \implies \mathbb{DC}(Z_G)$ is a face of $\mathbb{DC}(\Pi_n)$



Theorem (Padrol, Pilaud, P., '23)

Explicit facet-description of $\mathbb{DC}(Z_G)$

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Corollary

$\dim \mathbb{DC}(Z_G) = \# \text{ cliques of } G$

$\# \text{ facets of } \mathbb{DC}(Z_G) = \sum_{(i,j) \in E} 2^{|\{k ; (i,k), (j,k) \in E\}|}$

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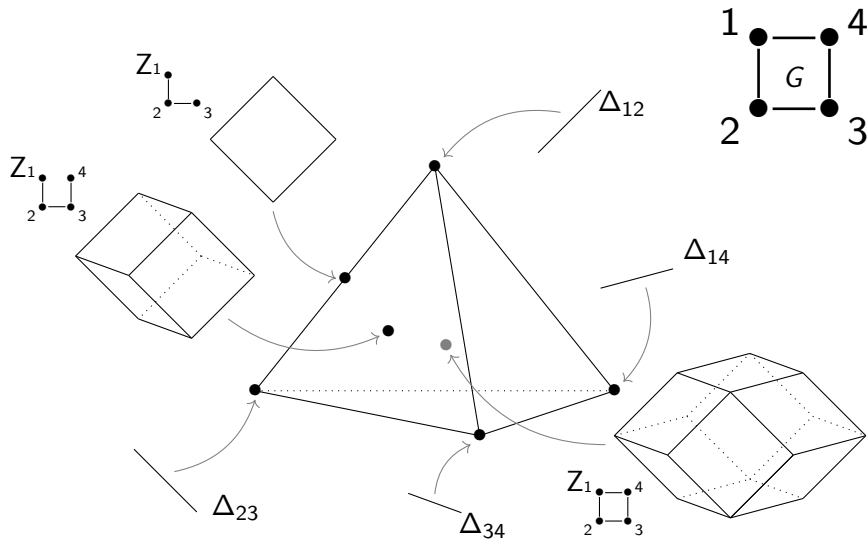
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Corollary

$\mathbb{DC}(Z_G)$ simplicial iff G without triangle

NB: Recover facet-description of $\mathbb{DC}(\Pi_n)$

My contribution - Graphical Zonotopes



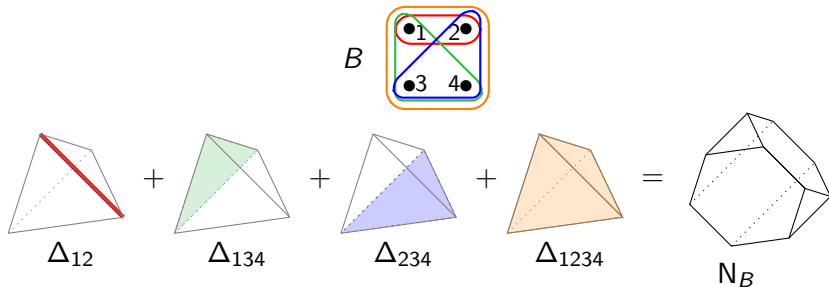
Definition

Building set $B \subseteq 2^{[n]}$ with: $X_{1,2} \in B, X_1 \cap X_2 \neq \emptyset \Rightarrow X_1 \cup X_2 \in B$

Definition

Nestohedron $N_B := \sum_{X \in B} \Delta_X$ where $\Delta_X = \text{conv}\{\mathbf{e}_i ; i \in X\}$

N_B deformation of $\Pi_n \Rightarrow \mathbb{DC}(N_B)$ is a face of $\mathbb{DC}(\Pi_n)$



Elementary blocks $X \in \varepsilon(B)$ iff X is not a union

Maximal block $\mu(X) := \max\{Y \in B ; Y \subsetneq X\}$

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$\dim \mathbb{DC}(N_B) = |B| - \# \text{ singletons}$

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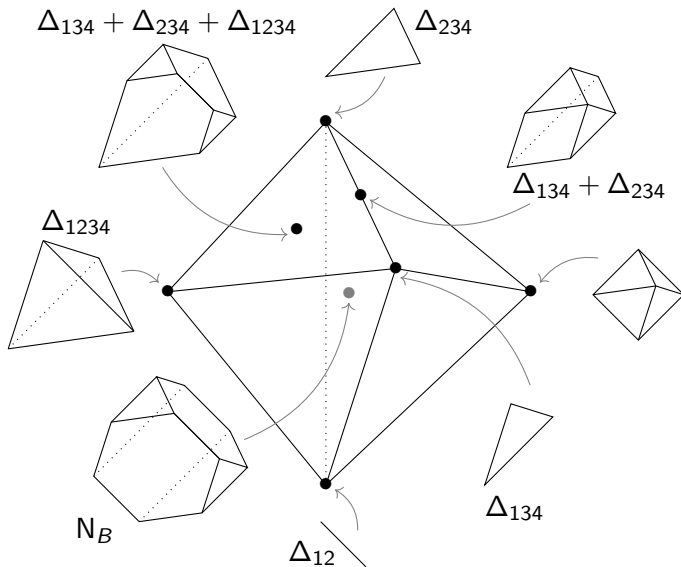
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Corollary

$\mathbb{DC}(N_B)$ simplicial iff B has no non-elementary block with 3 maximal subblocks

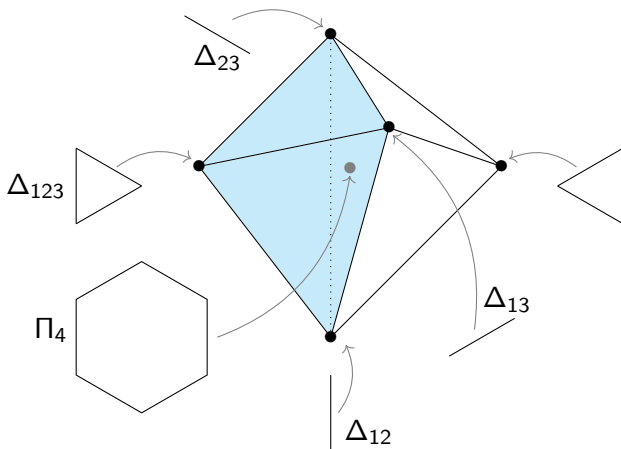
NB: Recover facet-description of $\mathbb{DC}(\Pi_n)$

My contribution - Nestohedra



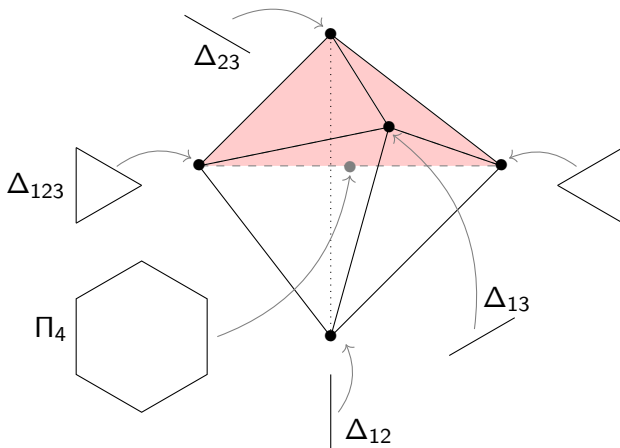
Definition

Hypergraphic pol $P_H := \sum_{X \in H} \Delta_X$ with $H \subseteq 2^{[n]}$



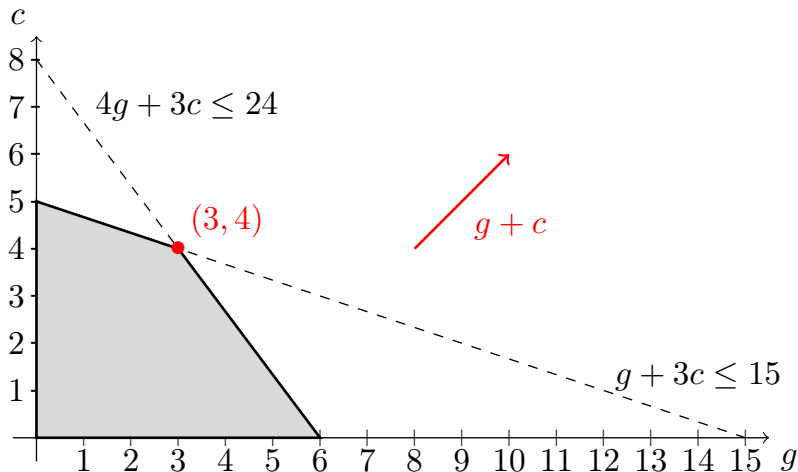
Definition

Quotientopes: Minkowski sum of shard polytopes

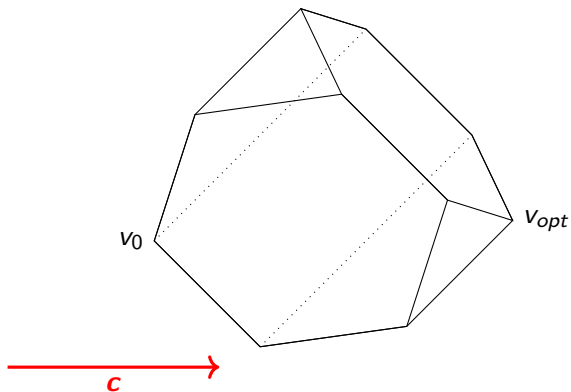


Max-slope Pivot Polytopes

Linear optimization

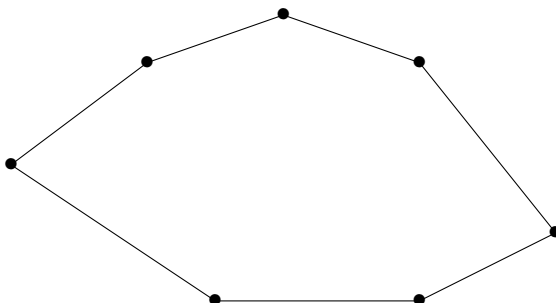


Simplex method



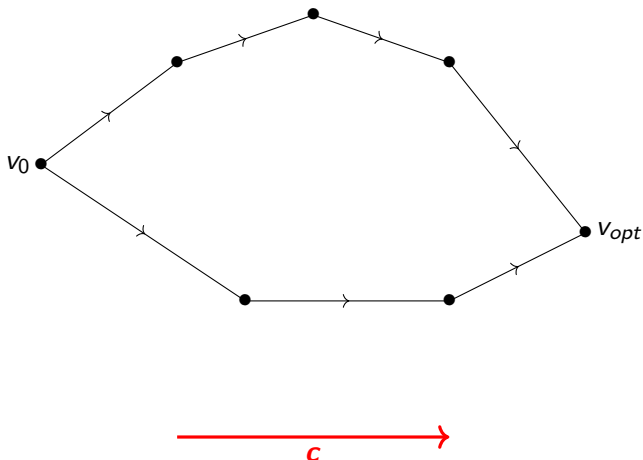
Max-slope pivot rule

Linear optimization in dimension 2 (simplex method):



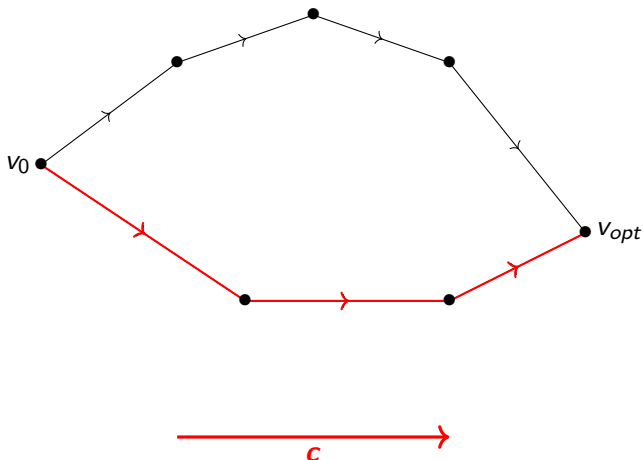
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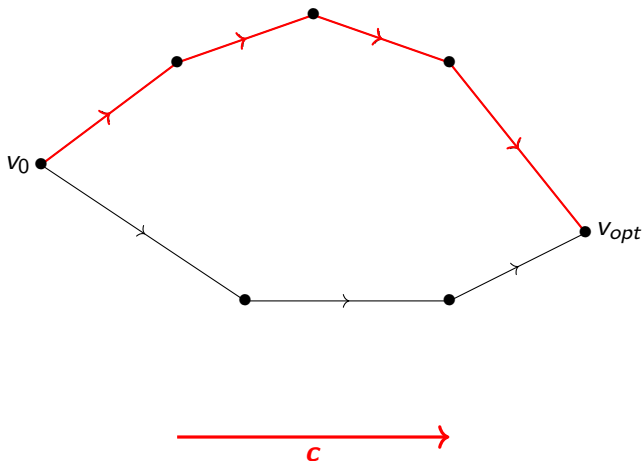
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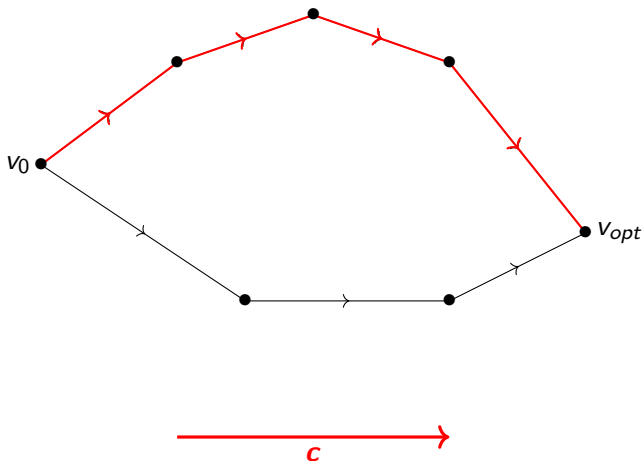
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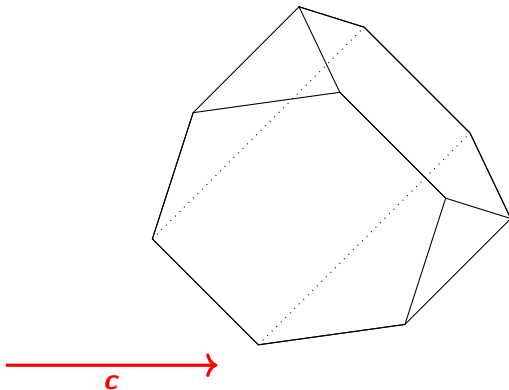
Linear optimization in dimension 2 (simplex method): **EASY !**



Convention: choose upper

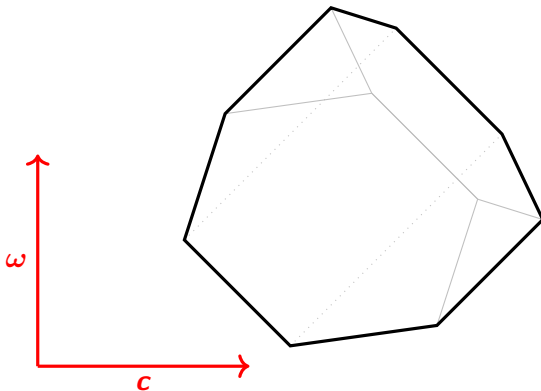
Max-slope pivot rule

Optimization in higher dimension: make it 2-dimensional !



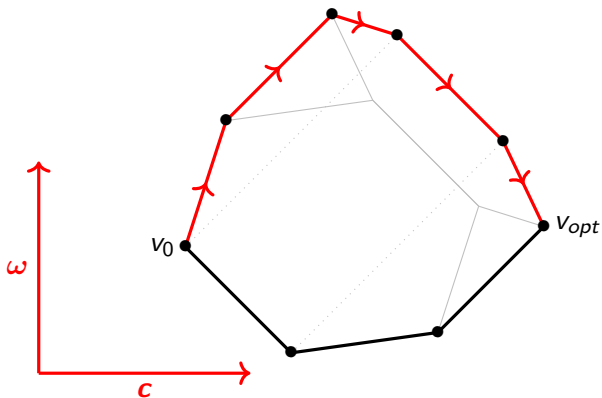
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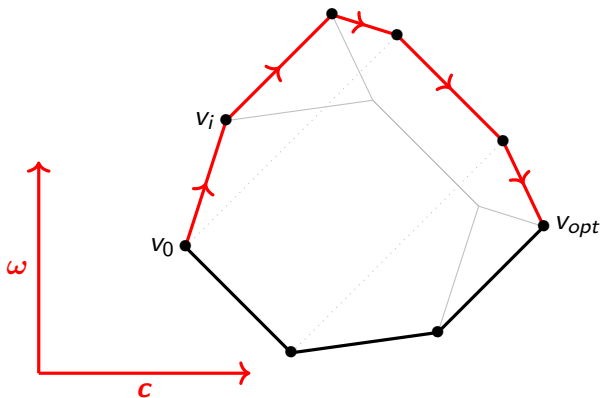
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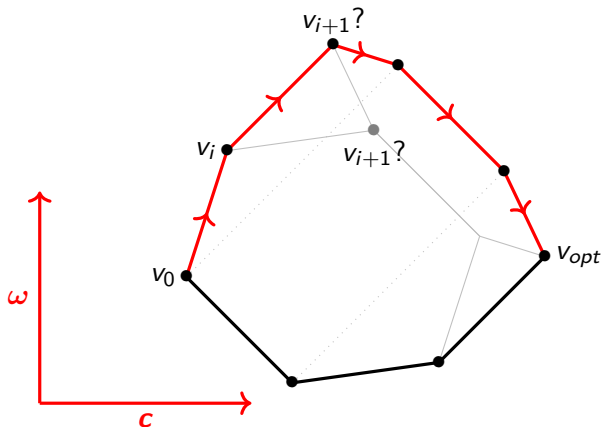
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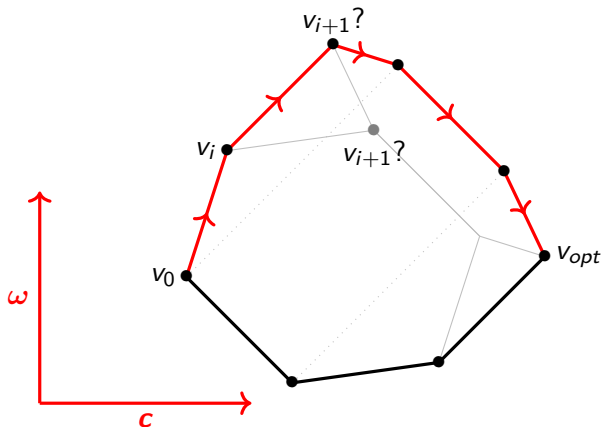
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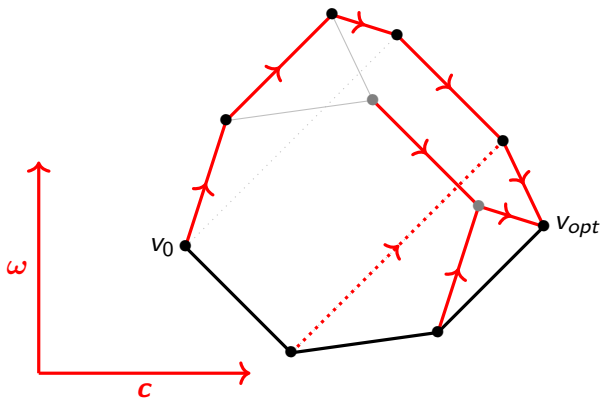
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Max-slope pivot rule: take (improving) neighbor with best slope

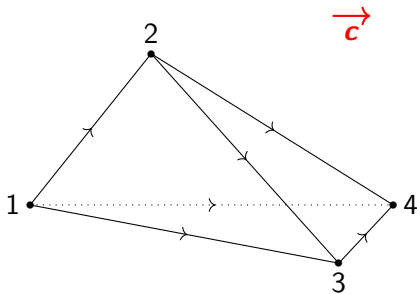
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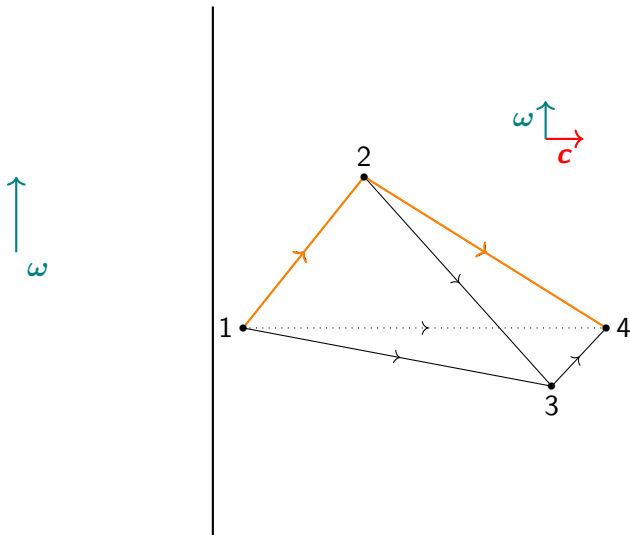


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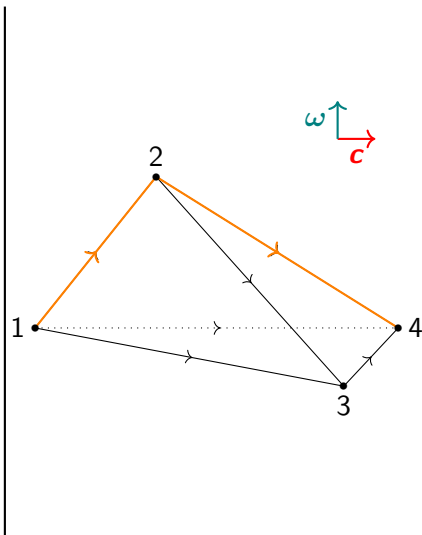
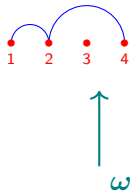
Coherent paths of the d -simplex



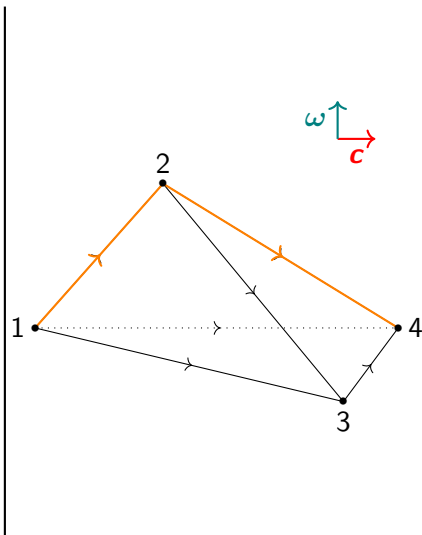
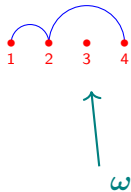
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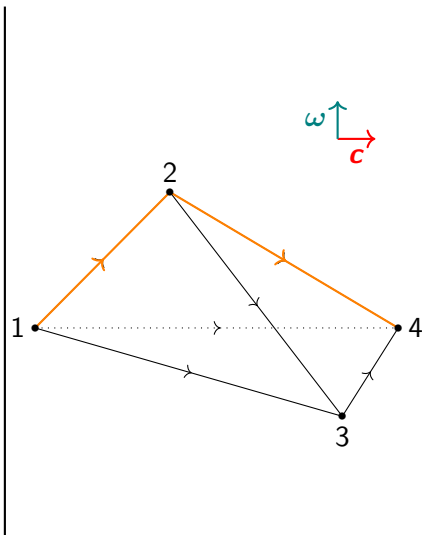
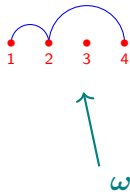
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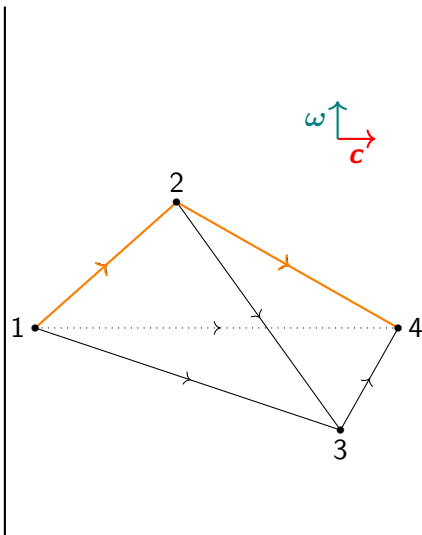
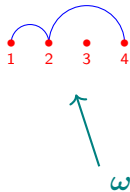
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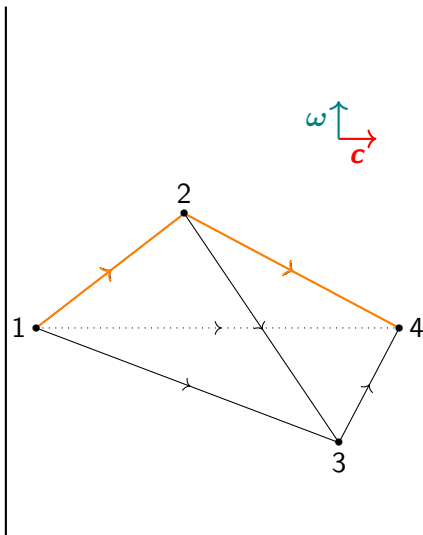
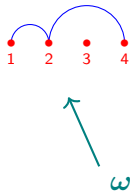
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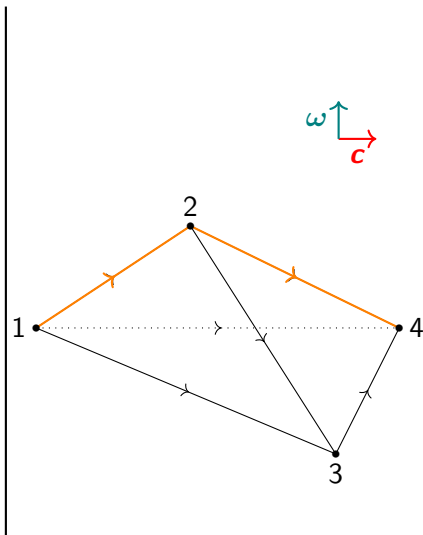
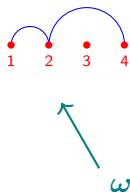
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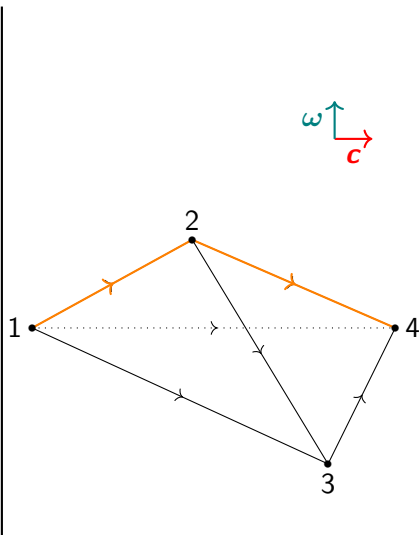
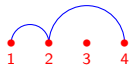
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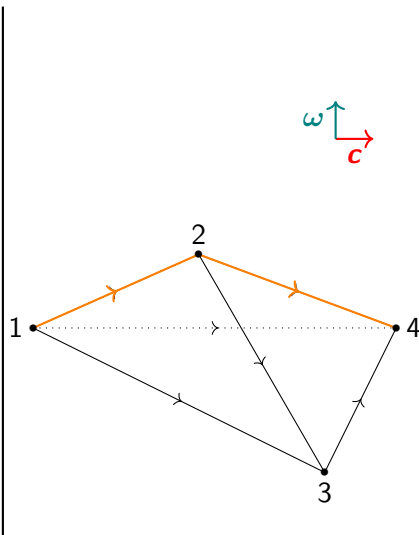
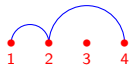
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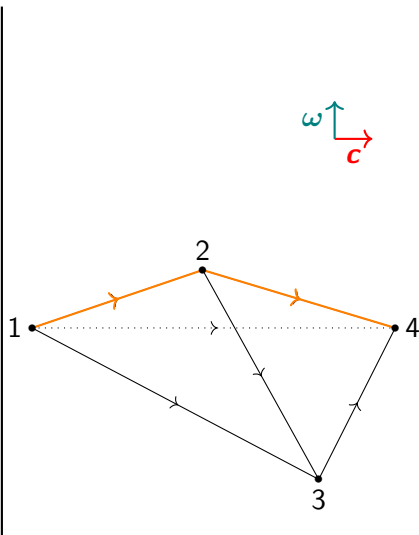
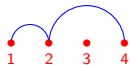
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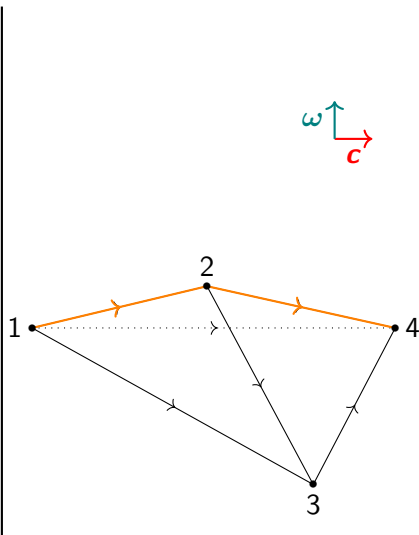
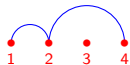
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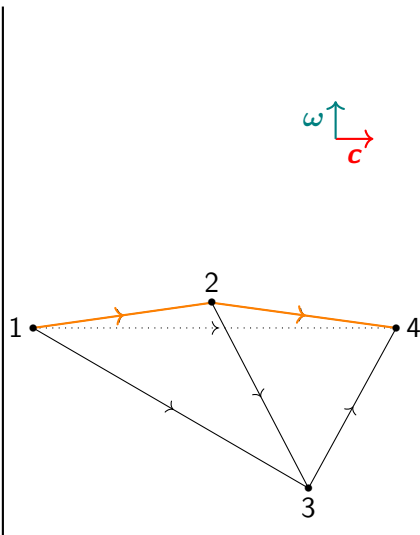
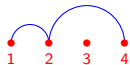
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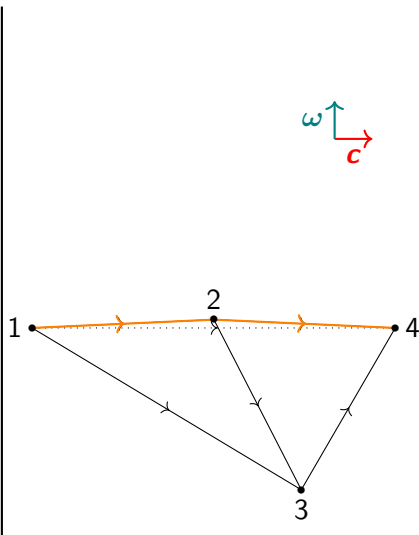
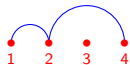
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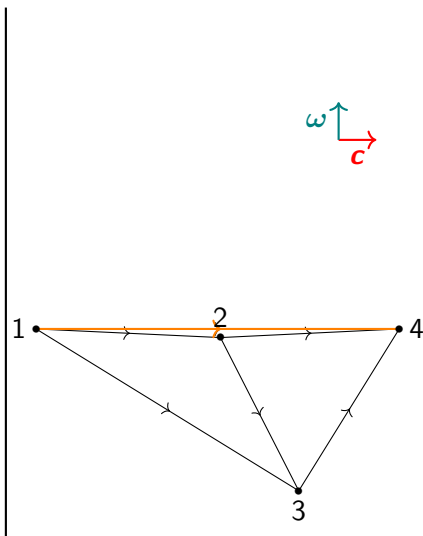
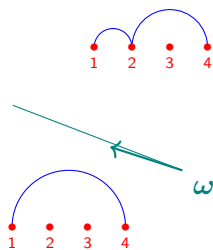
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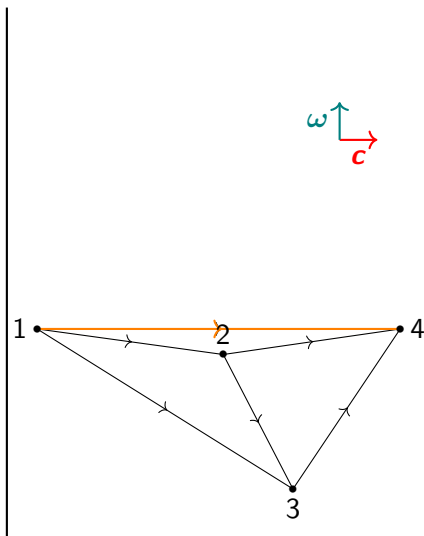
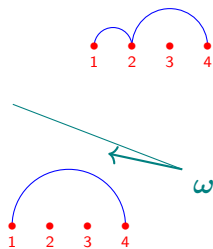
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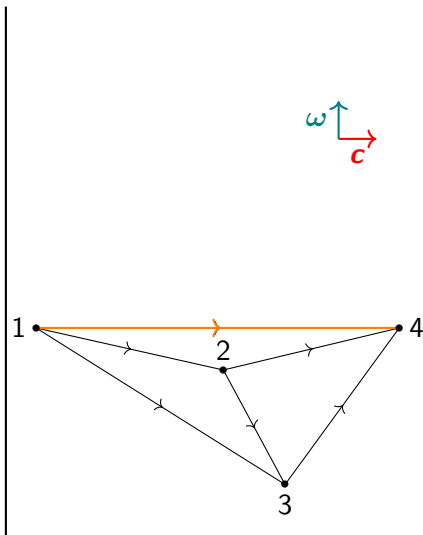
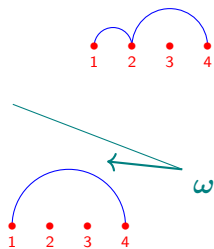
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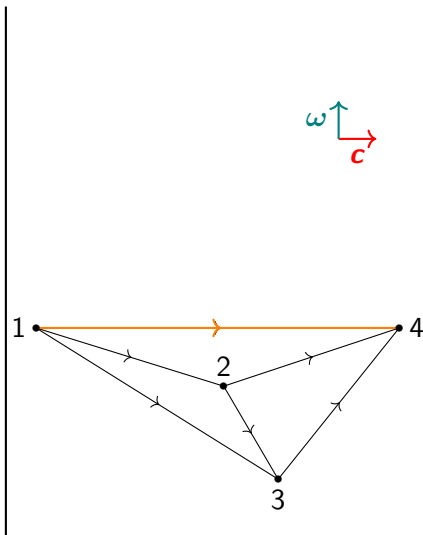
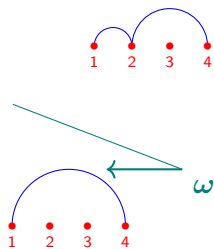
Coherent paths of the d -simplex



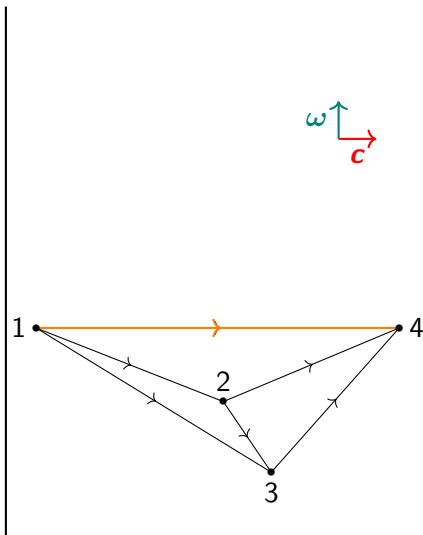
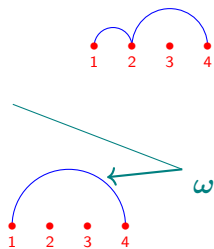
Coherent paths of the d -simplex



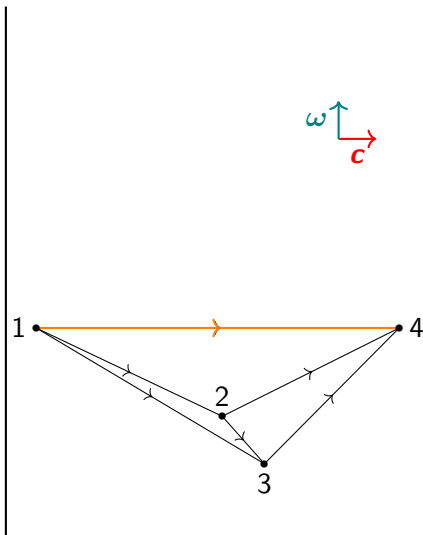
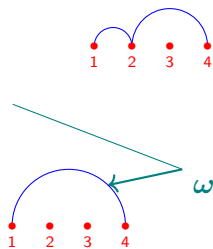
Coherent paths of the d -simplex



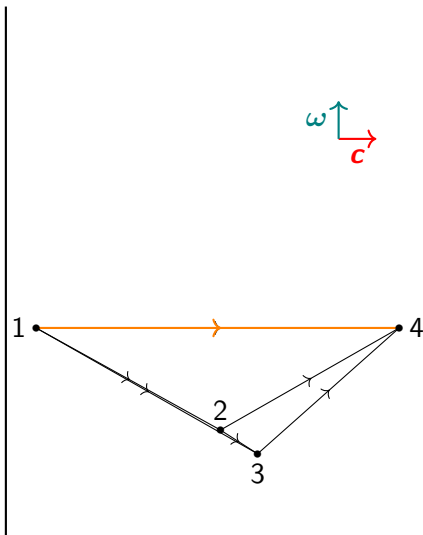
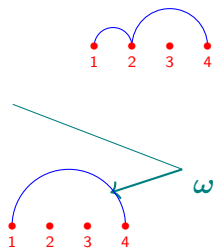
Coherent paths of the d -simplex



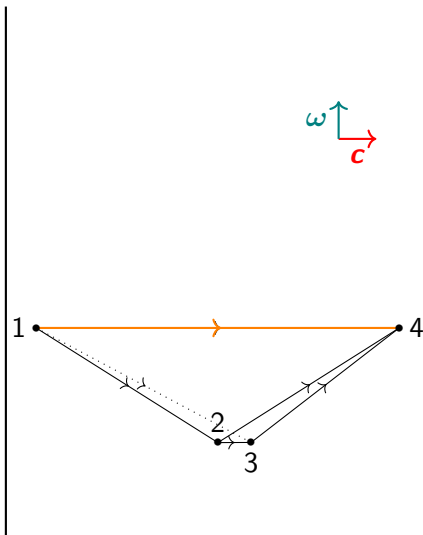
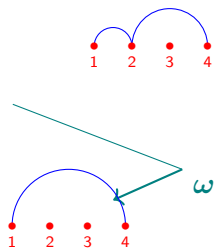
Coherent paths of the d -simplex



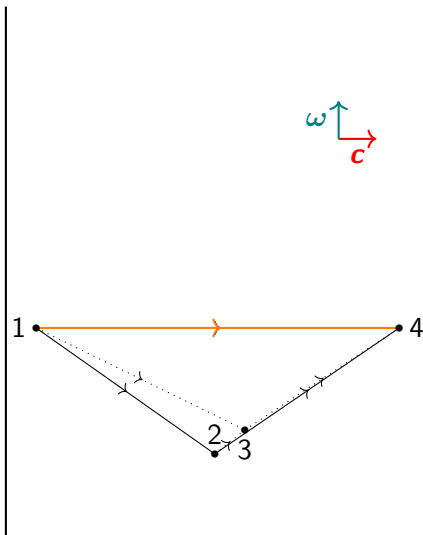
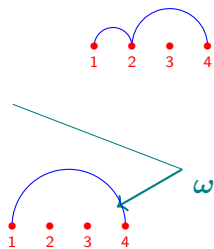
Coherent paths of the d -simplex



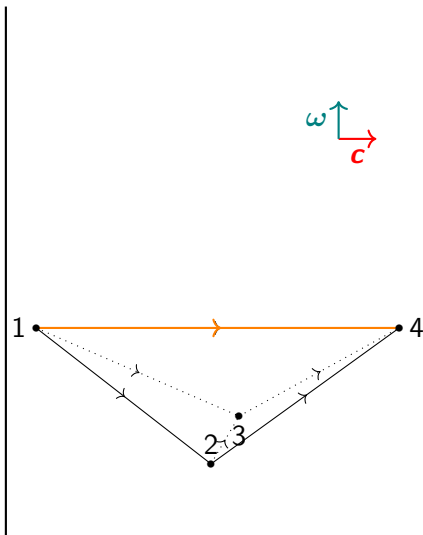
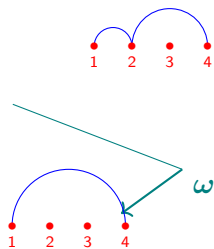
Coherent paths of the d -simplex



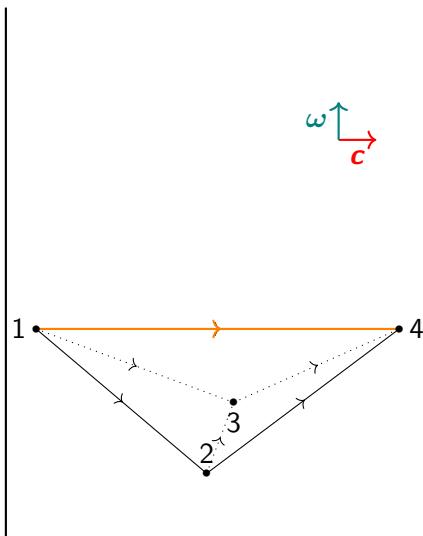
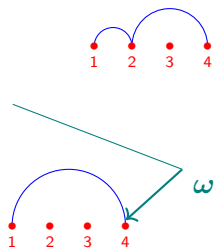
Coherent paths of the d -simplex



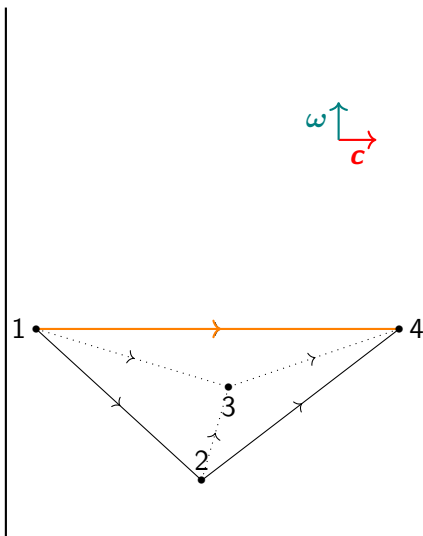
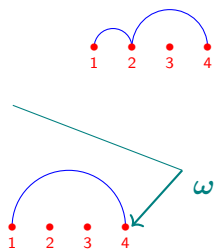
Coherent paths of the d -simplex



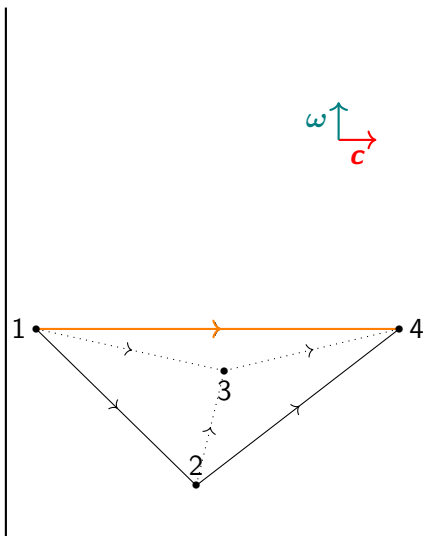
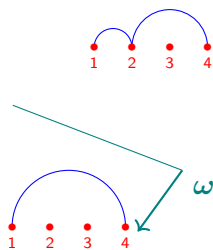
Coherent paths of the d -simplex



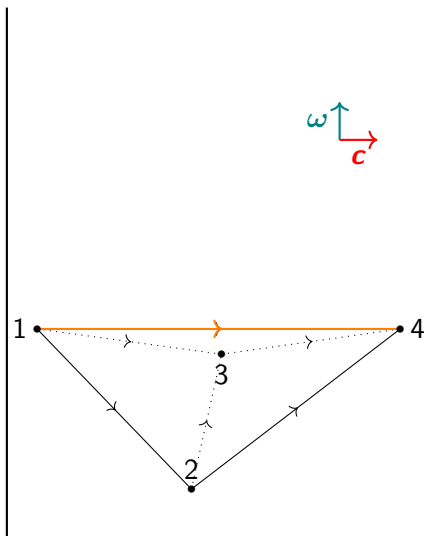
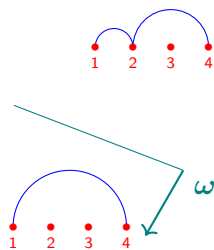
Coherent paths of the d -simplex



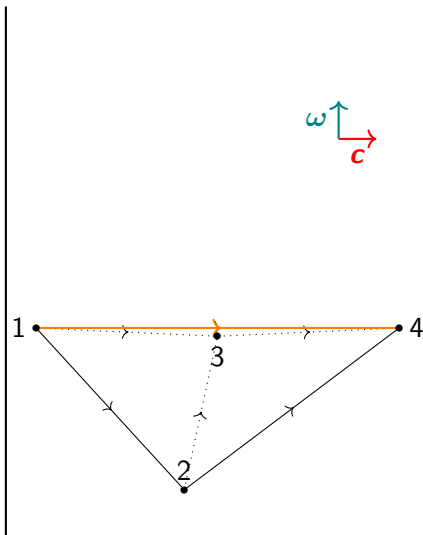
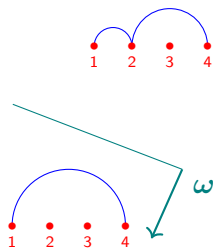
Coherent paths of the d -simplex



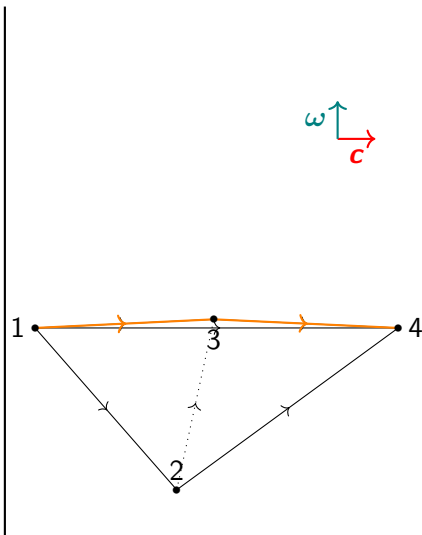
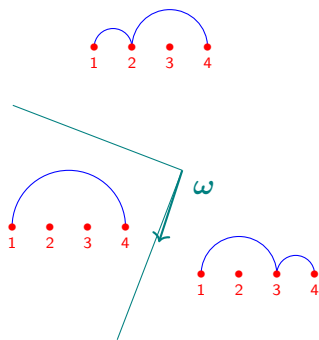
Coherent paths of the d -simplex



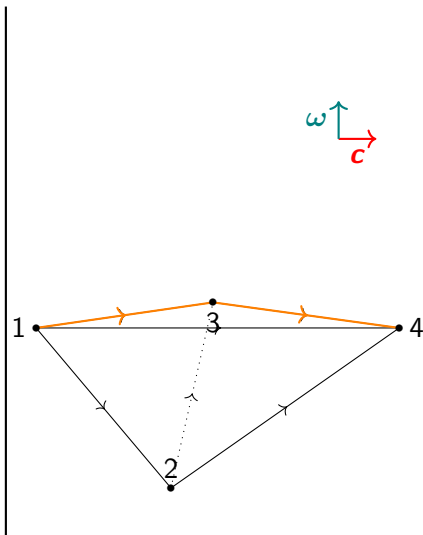
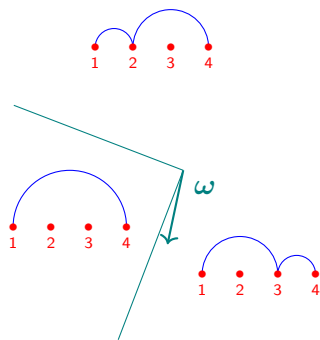
Coherent paths of the d -simplex



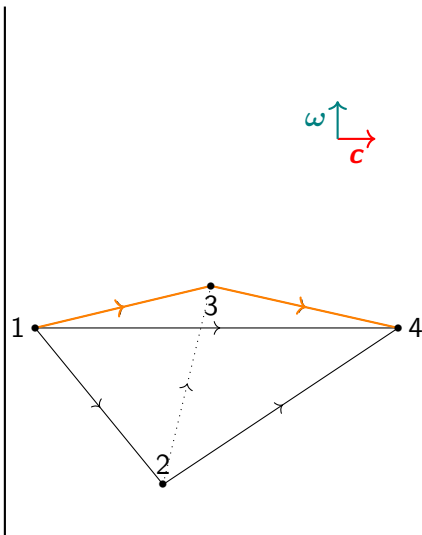
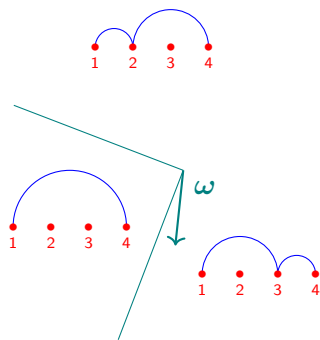
Coherent paths of the d -simplex



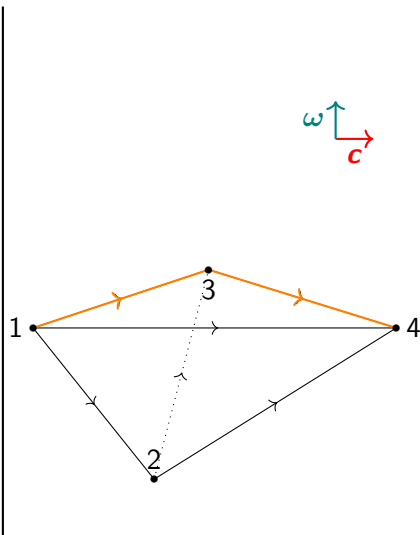
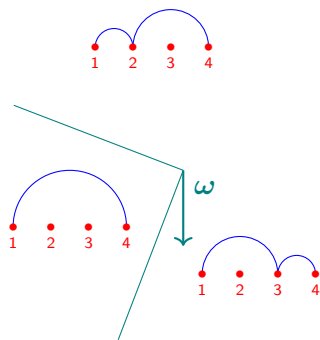
Coherent paths of the d -simplex



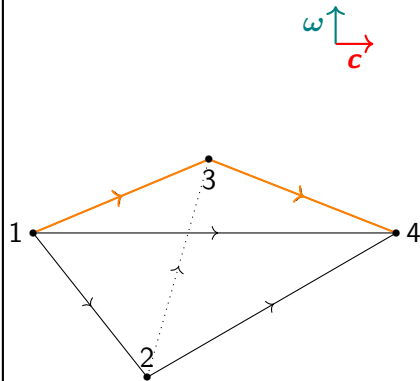
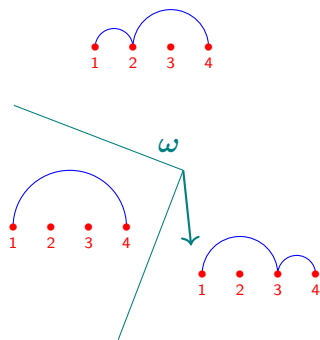
Coherent paths of the d -simplex



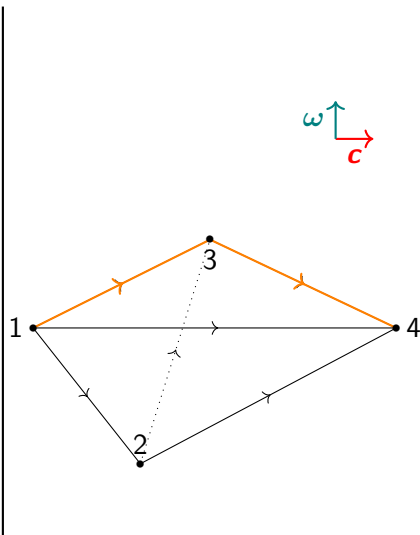
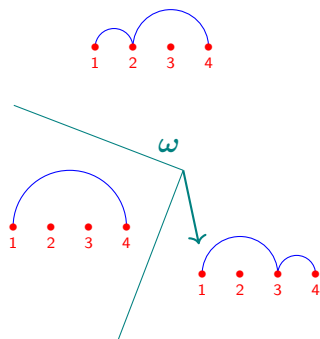
Coherent paths of the d -simplex



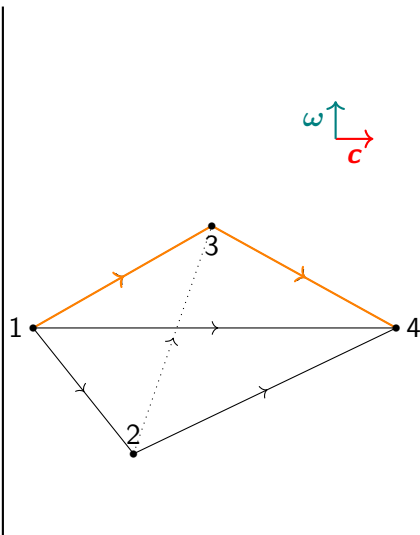
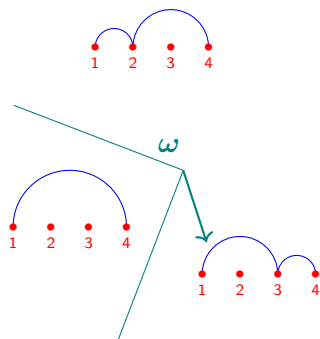
Coherent paths of the d -simplex



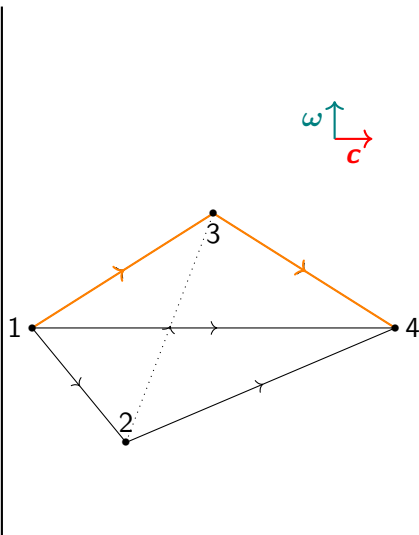
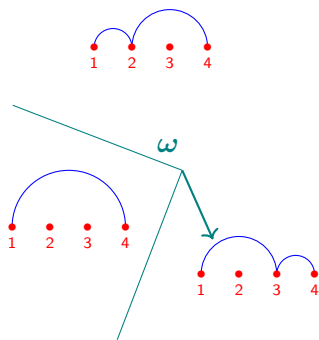
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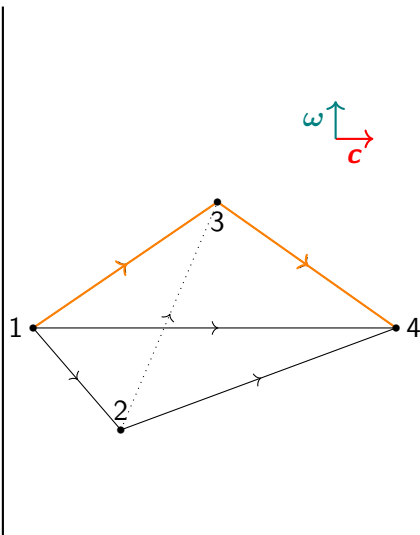
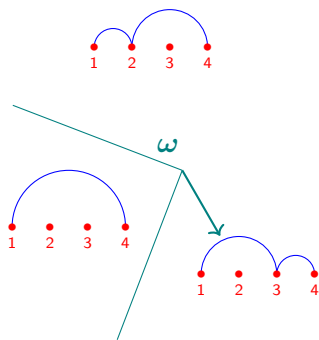
Coherent paths of the d -simplex



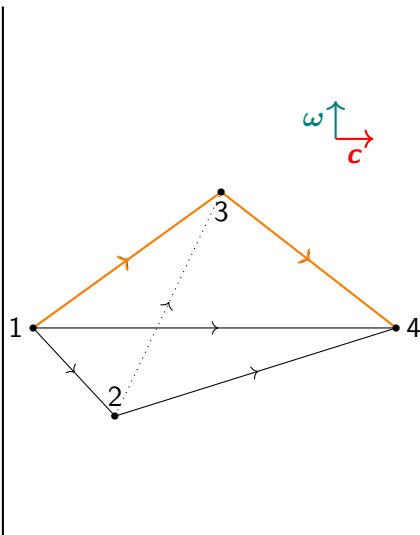
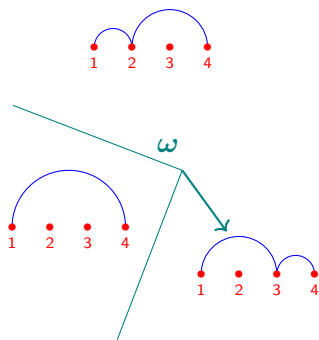
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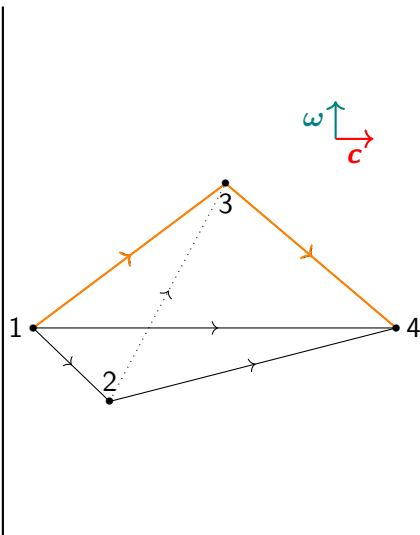
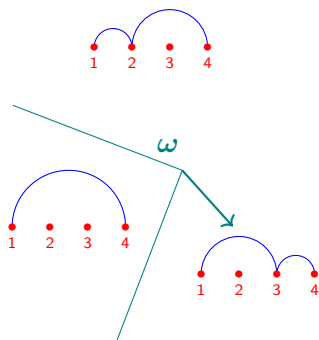
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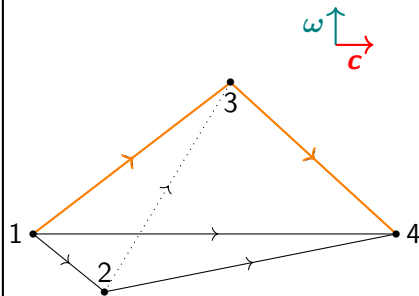
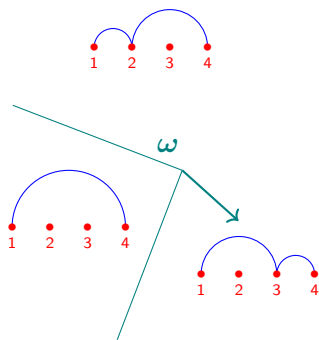
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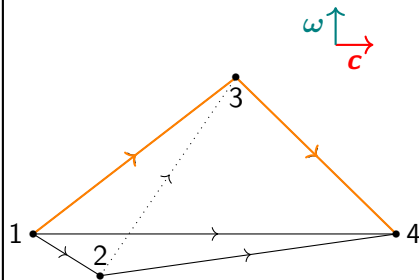
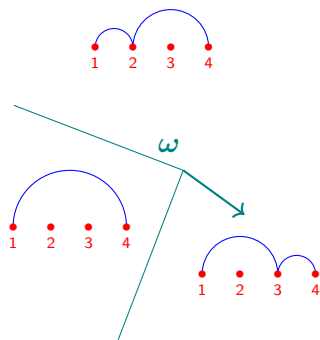
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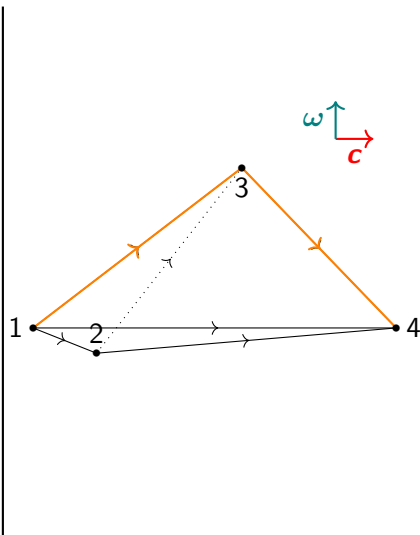
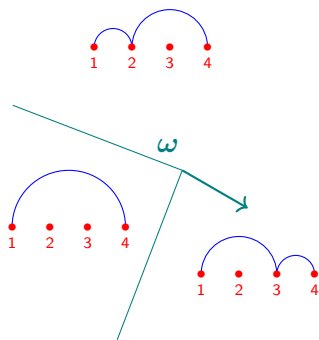
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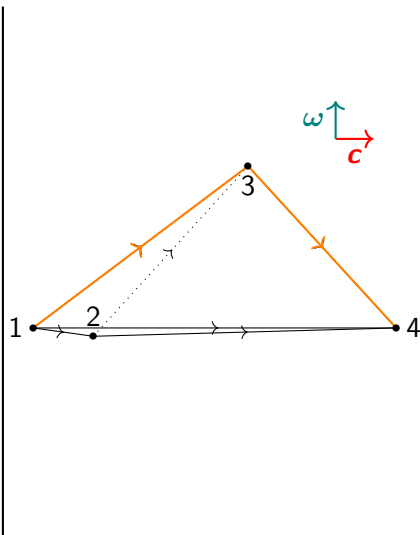
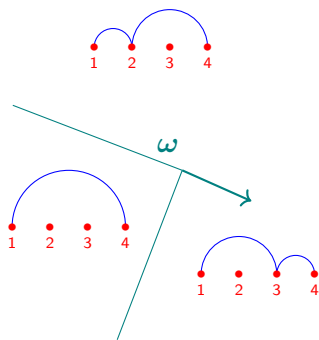
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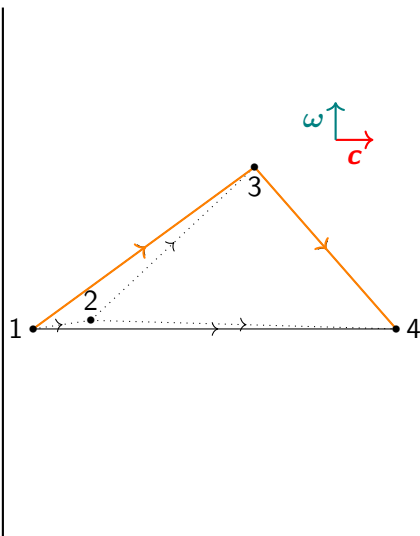
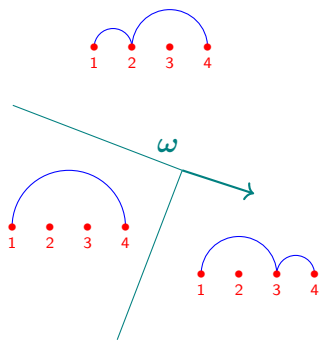
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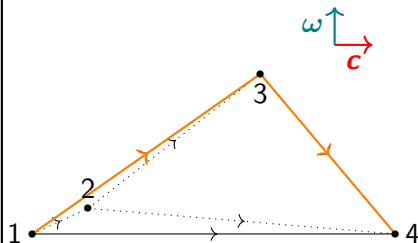
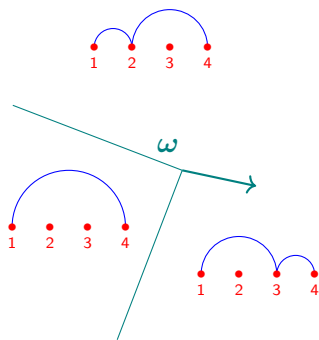
Coherent paths of the d -simplex



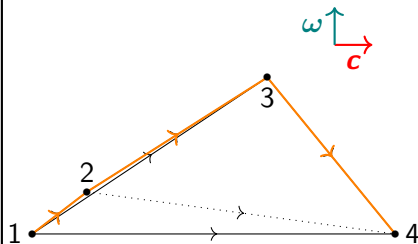
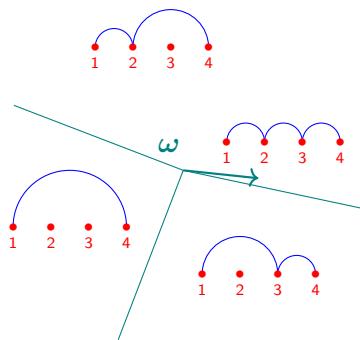
Coherent paths of the d -simplex



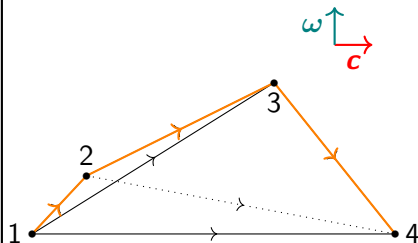
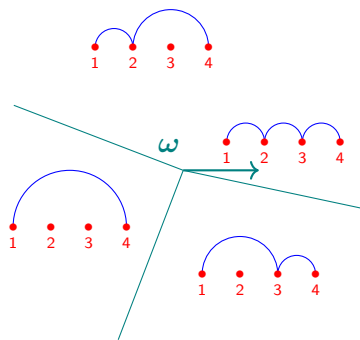
Coherent paths of the d -simplex



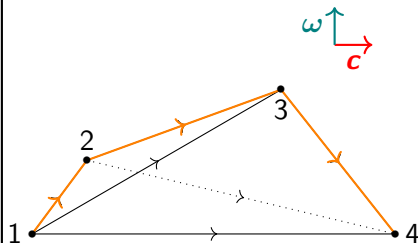
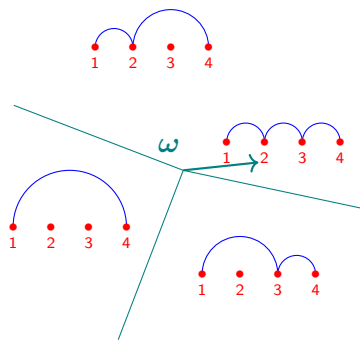
Coherent paths of the d -simplex



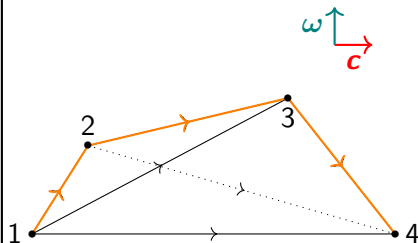
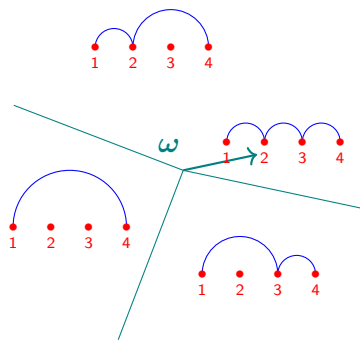
Coherent paths of the d -simplex



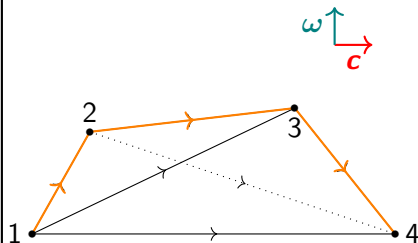
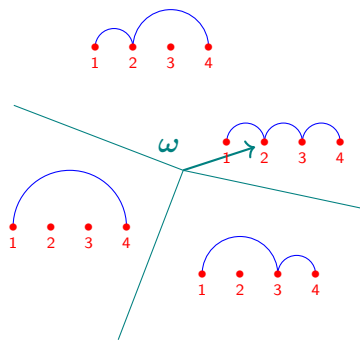
Coherent paths of the d -simplex



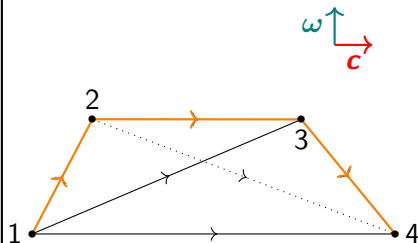
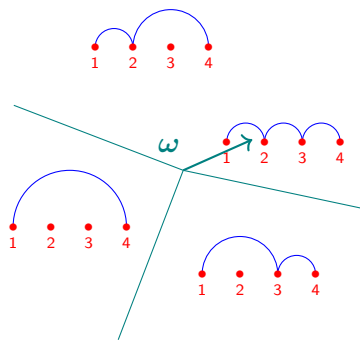
Coherent paths of the d -simplex



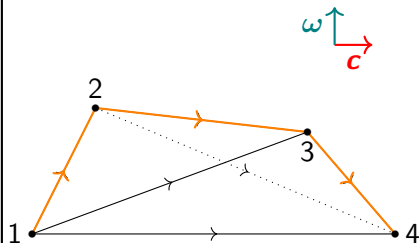
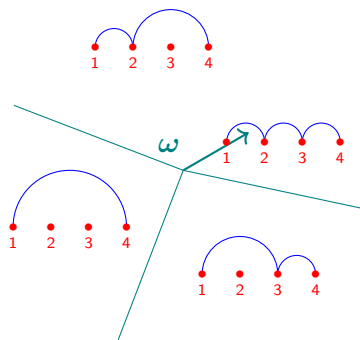
Coherent paths of the d -simplex



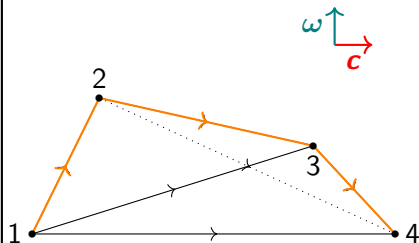
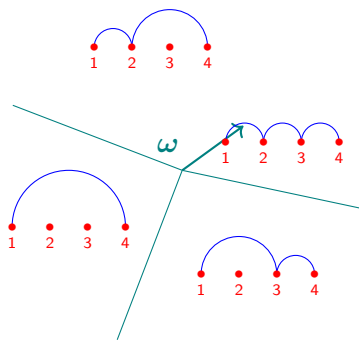
Coherent paths of the d -simplex



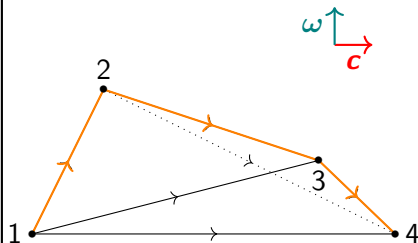
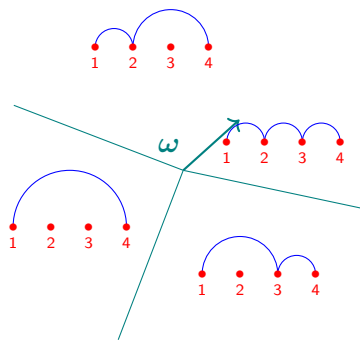
Coherent paths of the d -simplex



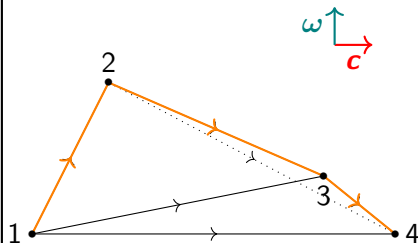
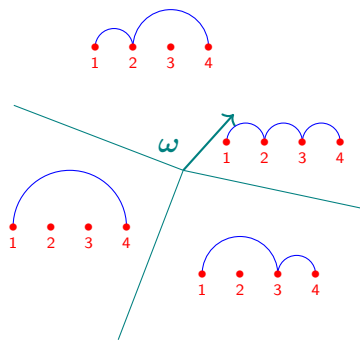
Coherent paths of the d -simplex



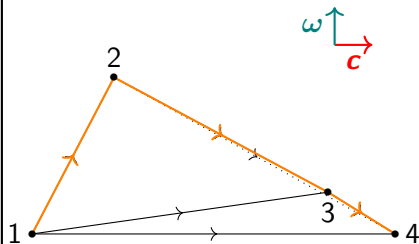
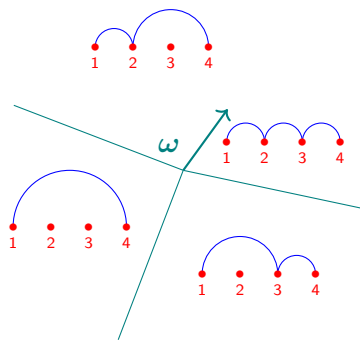
Coherent paths of the d -simplex



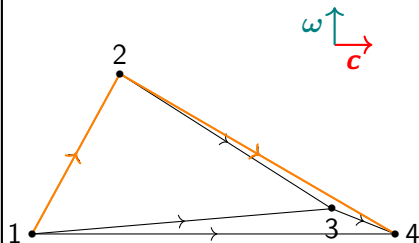
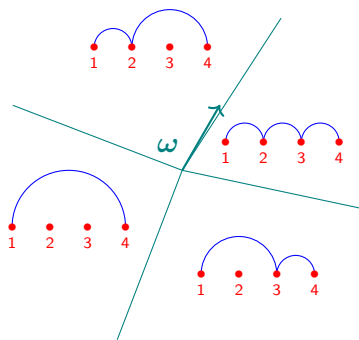
Coherent paths of the d -simplex



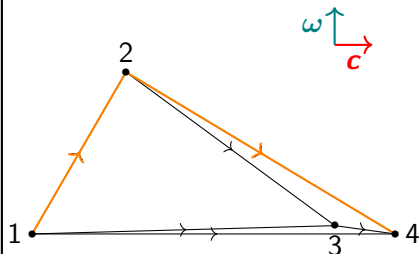
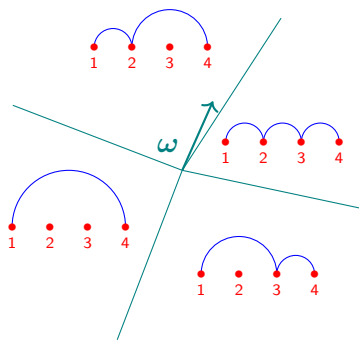
Coherent paths of the d -simplex



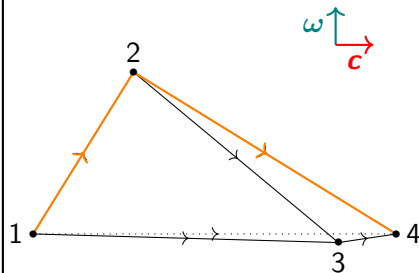
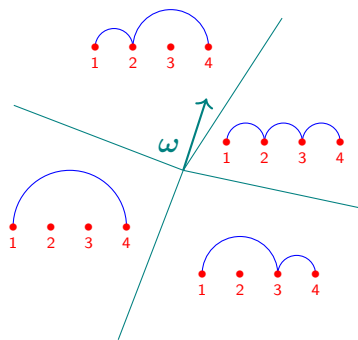
Coherent paths of the d -simplex



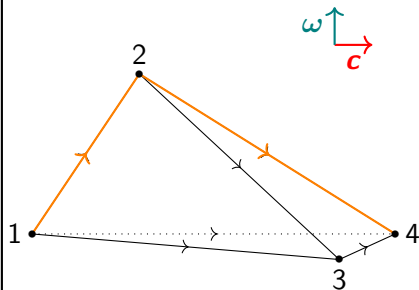
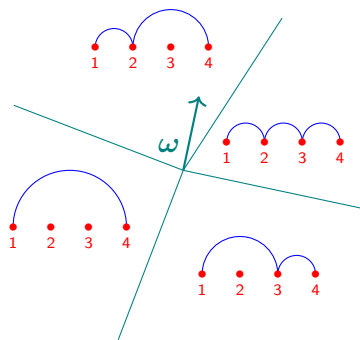
Coherent paths of the d -simplex



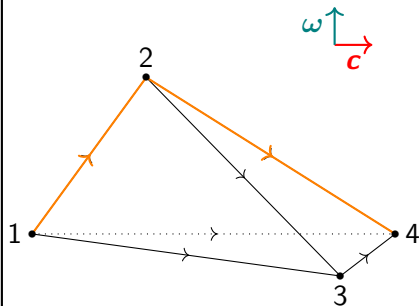
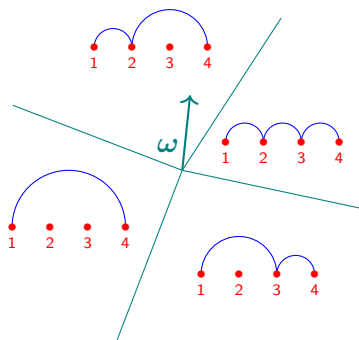
Coherent paths of the d -simplex



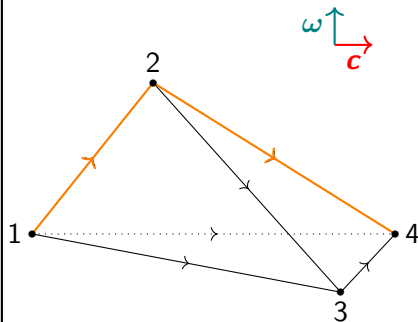
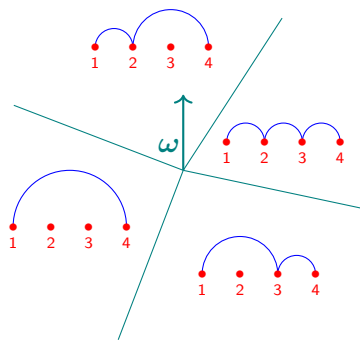
Coherent paths of the d -simplex



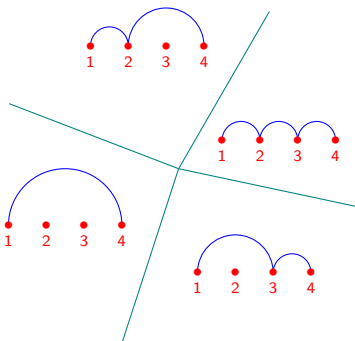
Coherent paths of the d -simplex



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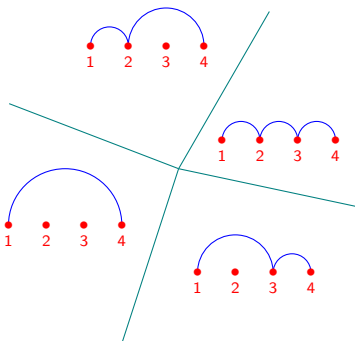
Monotone path polytope



Coherent monotone path: path obtained via max-slope pivot rule

Monotone path fan: Fan with $\omega \sim \omega'$ iff same path

Monotone path polytope



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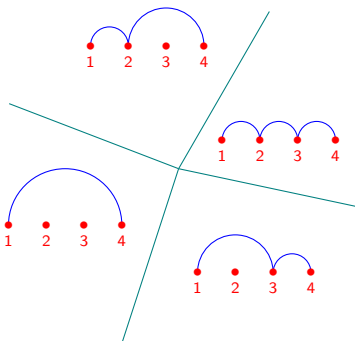
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Theorem (Billera, Sturmfels, '92)

The monotone path fan is polytopal.

Monotone path polytope $\Sigma_c(P)$: dual to monotone path fan

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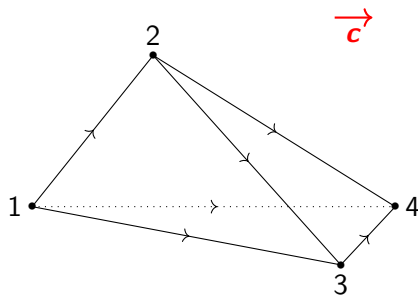
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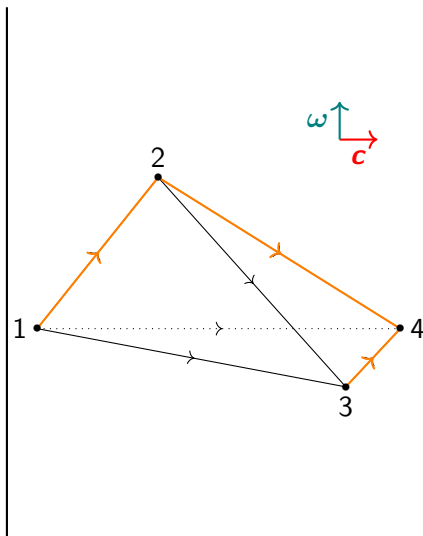
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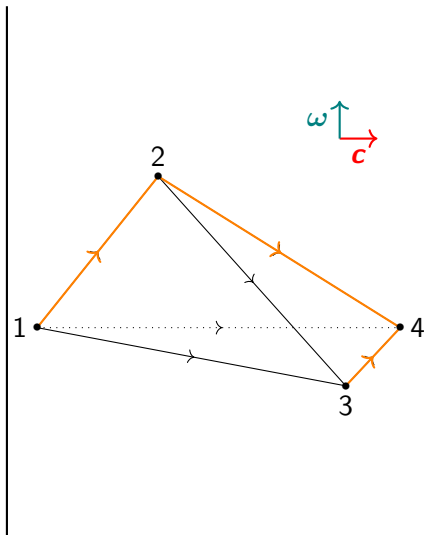
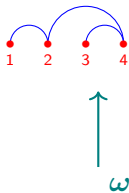
Coherent arborescences of the d -simplex



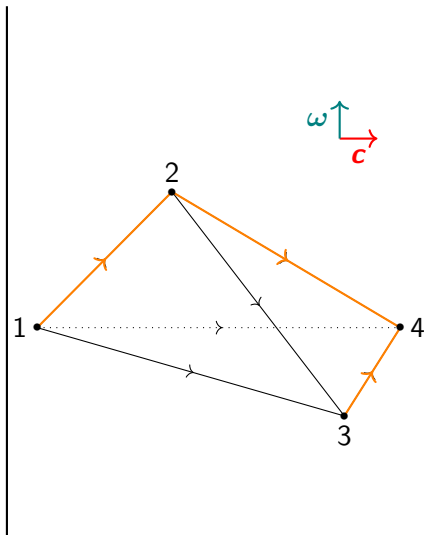
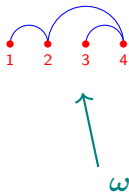
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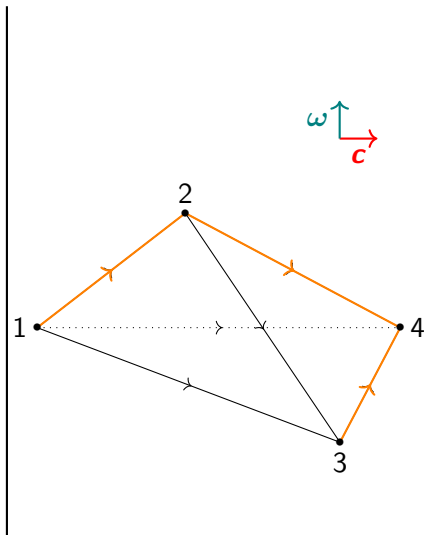
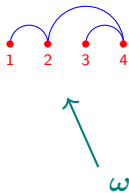
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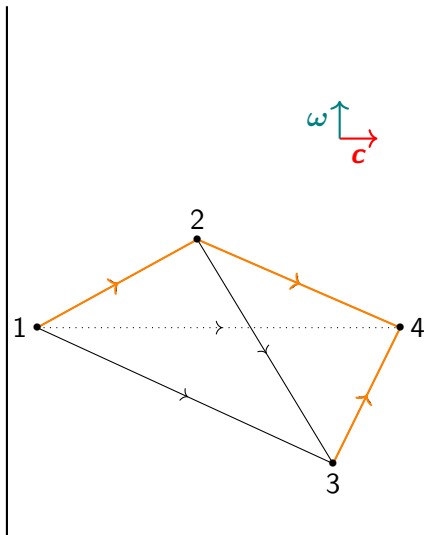
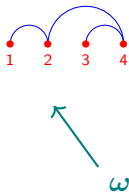
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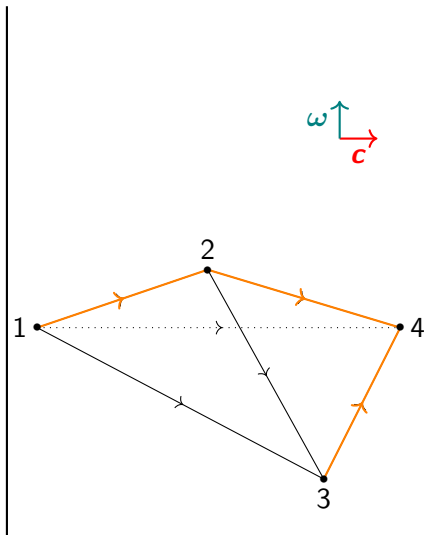
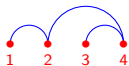
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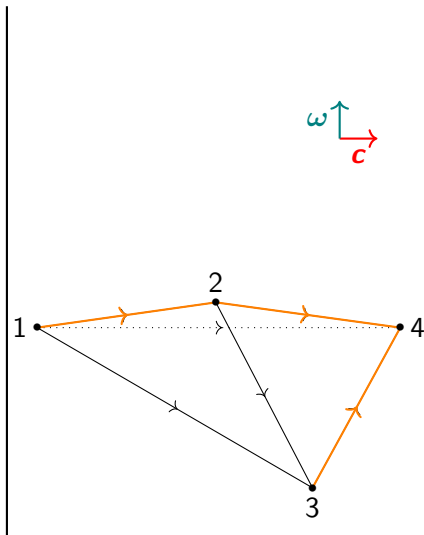
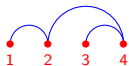
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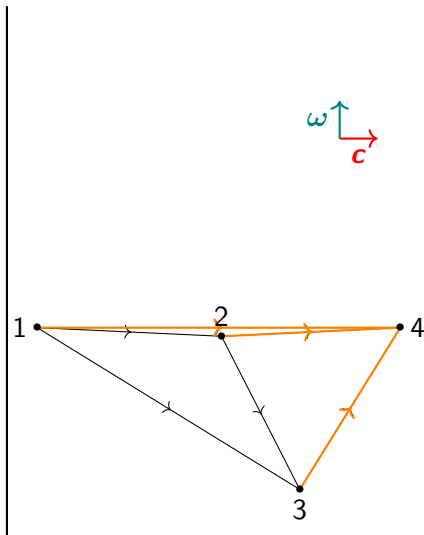
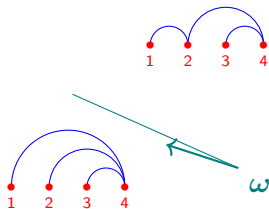
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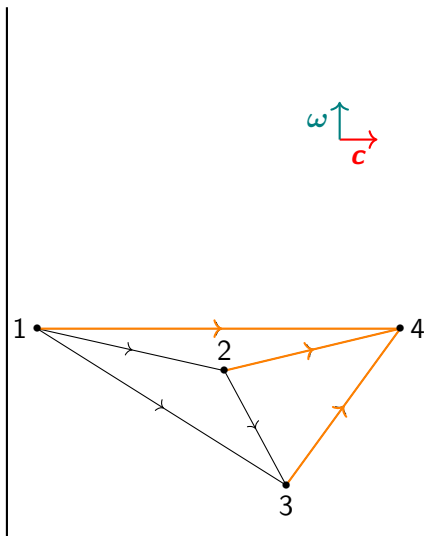
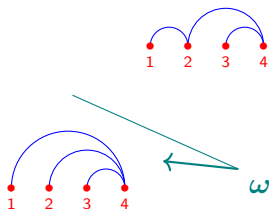
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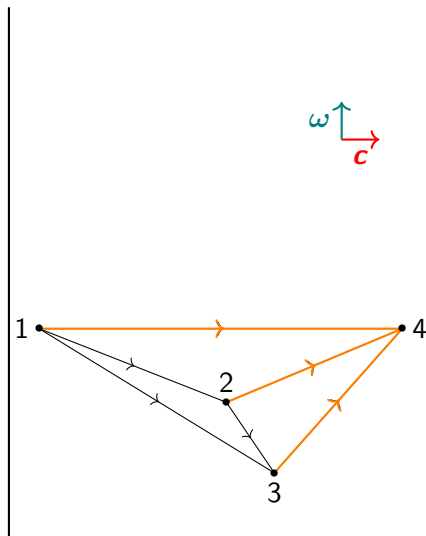
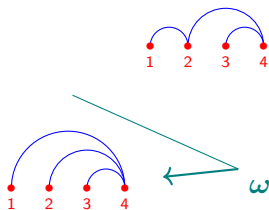
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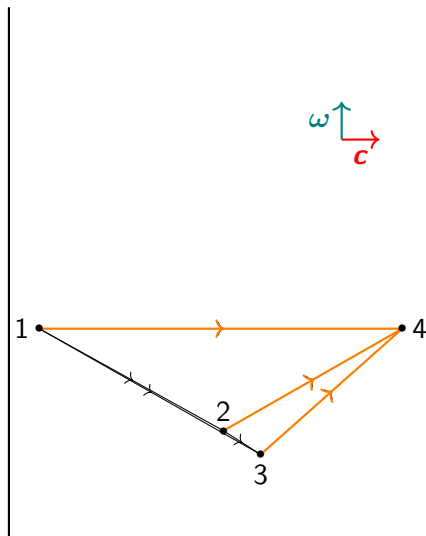
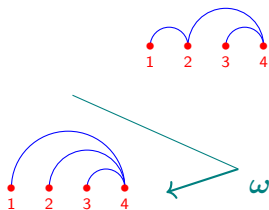
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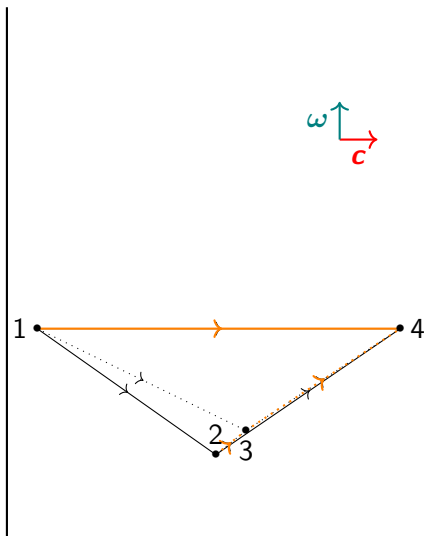
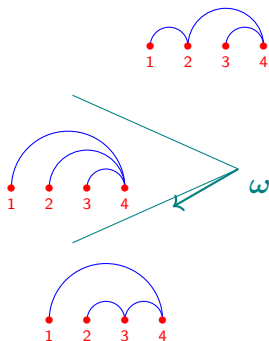
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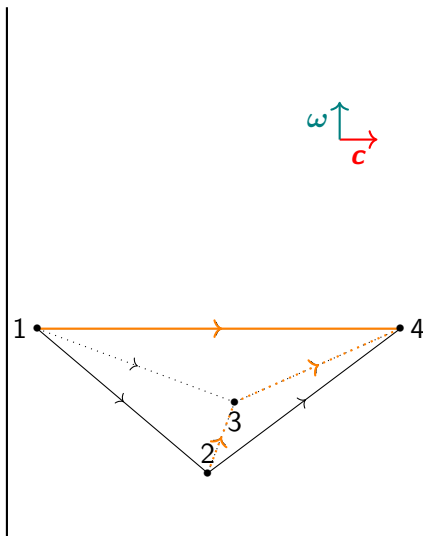
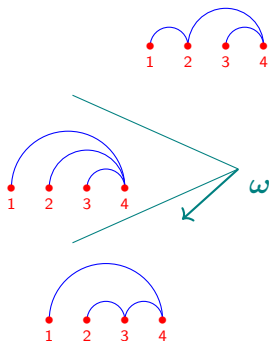
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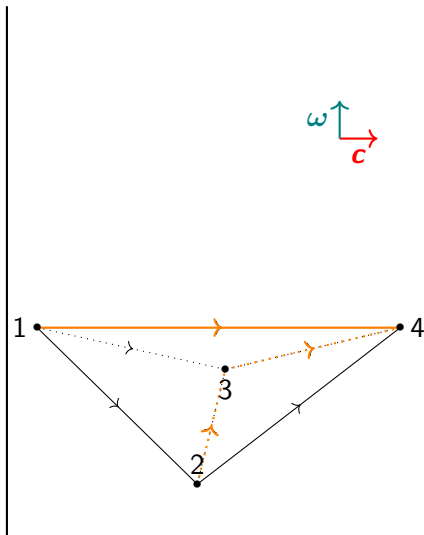
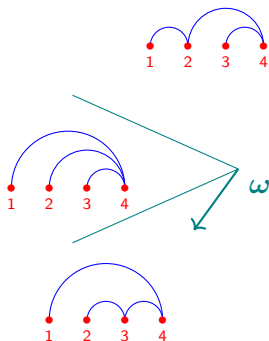
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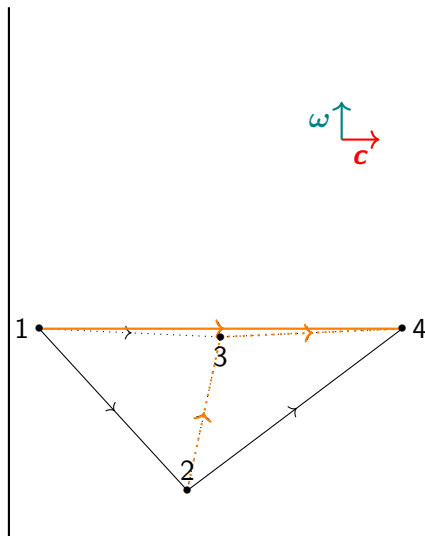
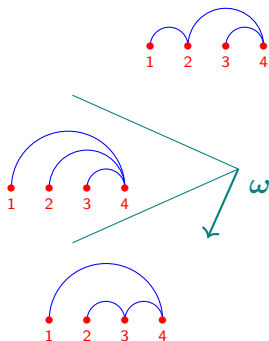
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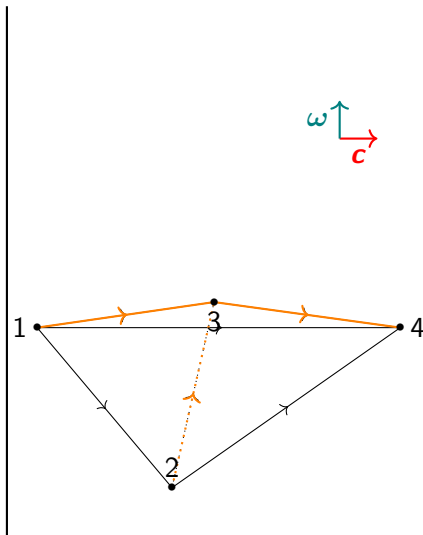
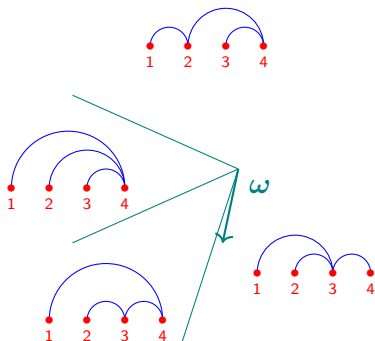
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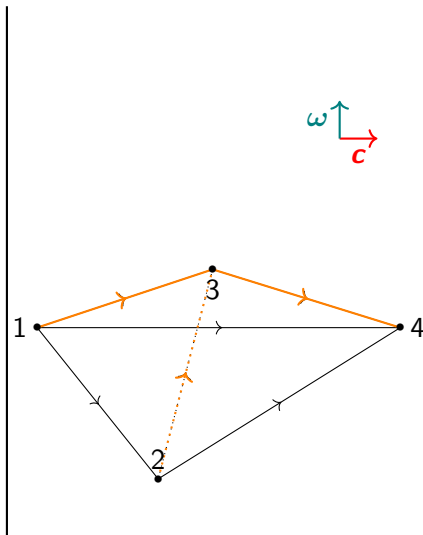
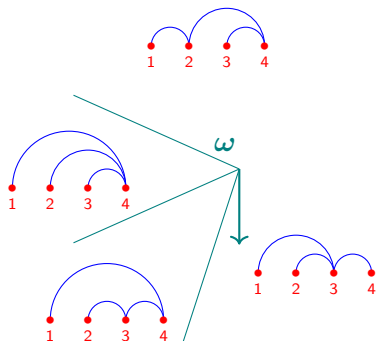
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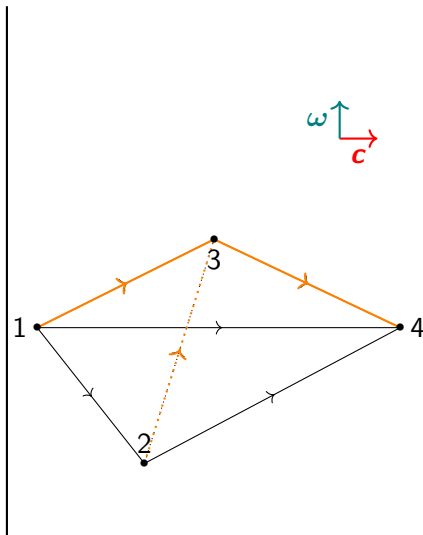
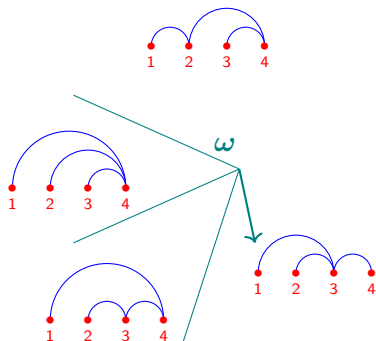
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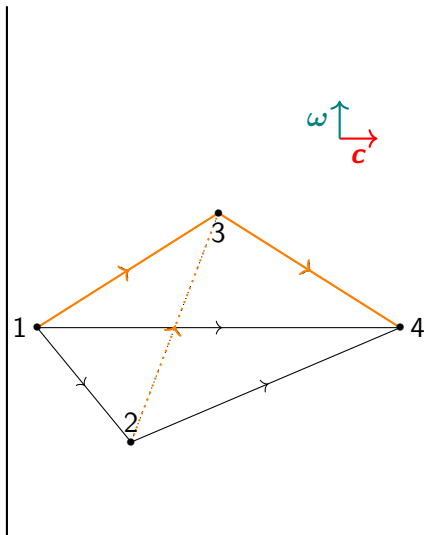
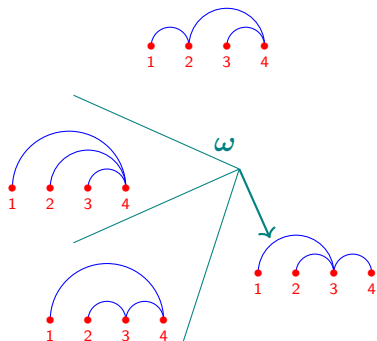
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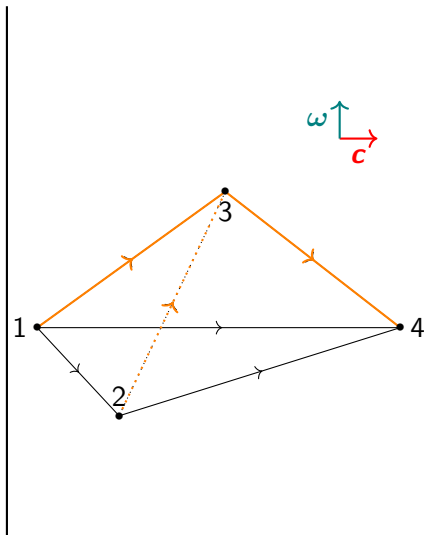
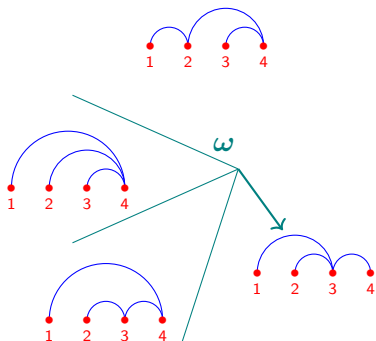
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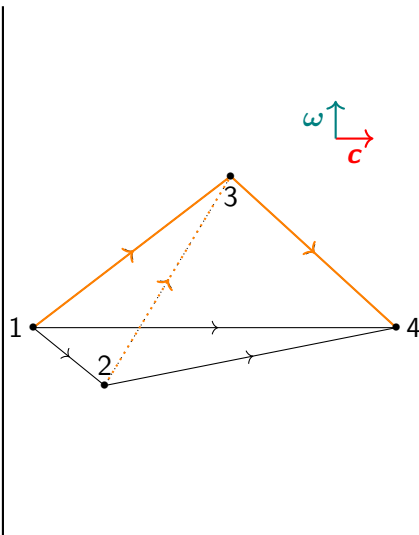
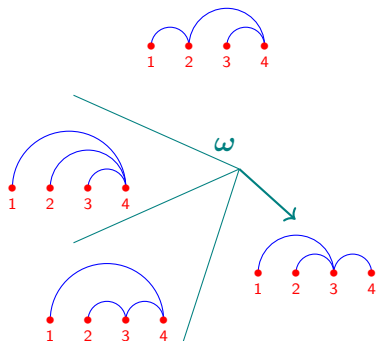
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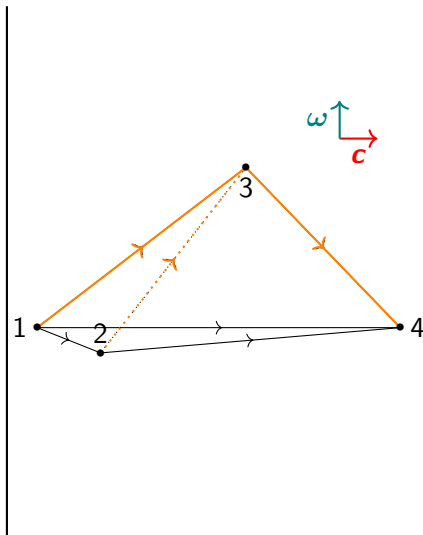
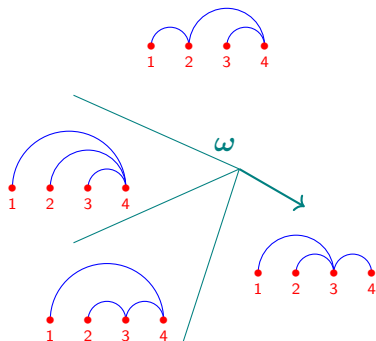
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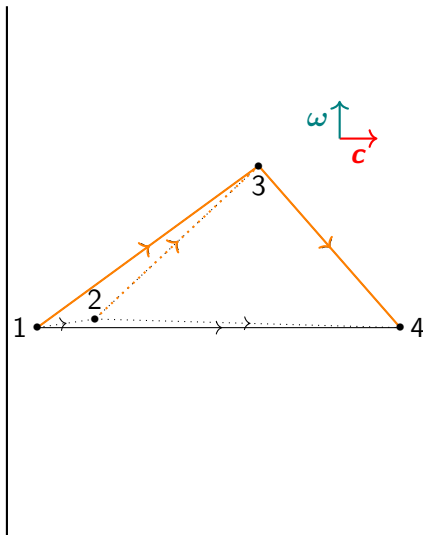
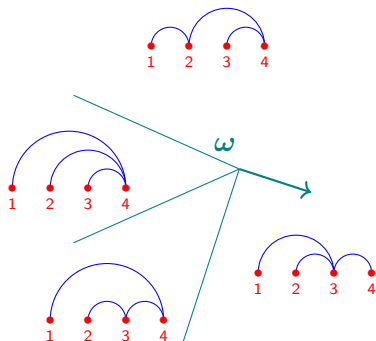
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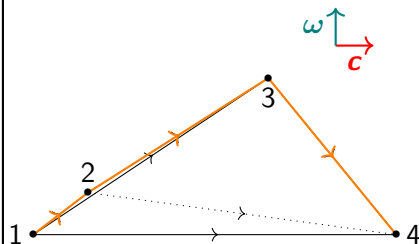
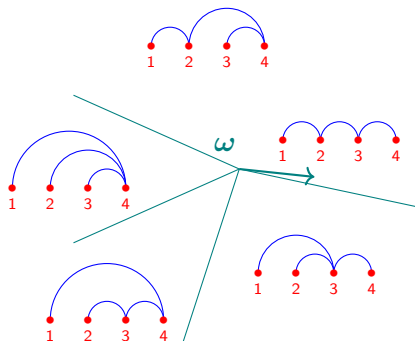
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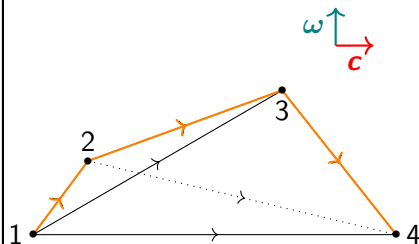
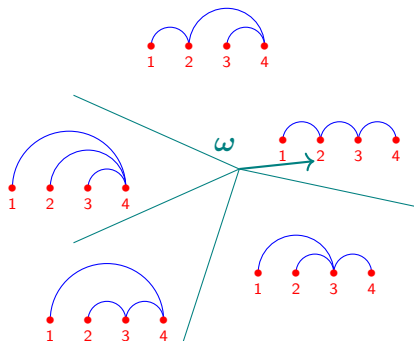
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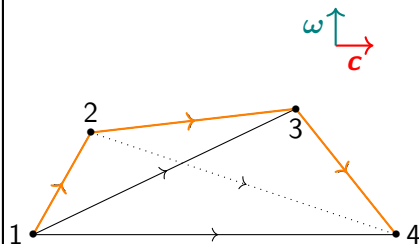
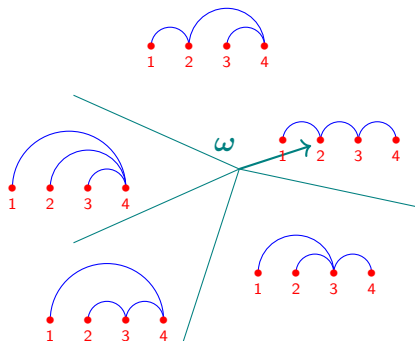
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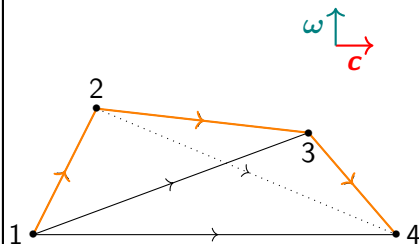
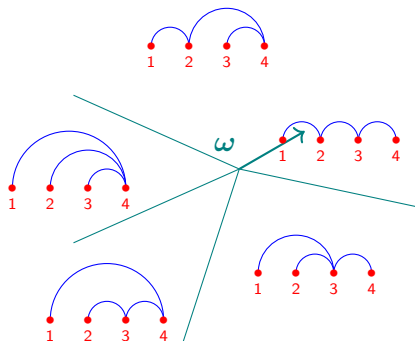
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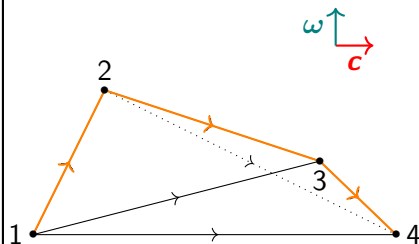
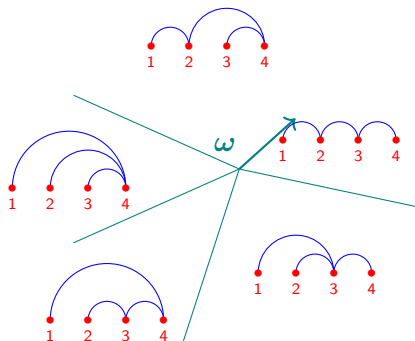
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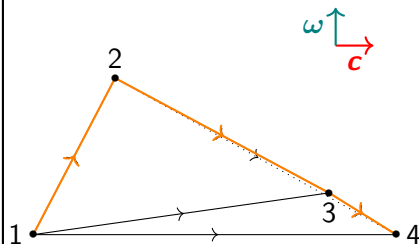
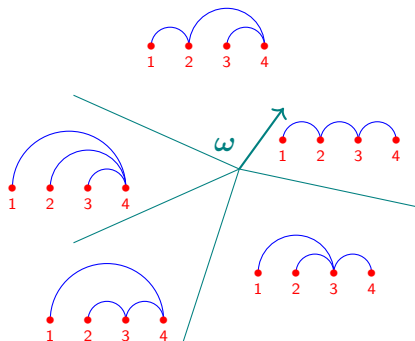
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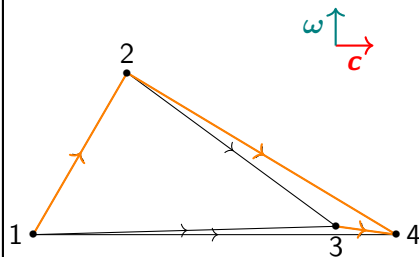
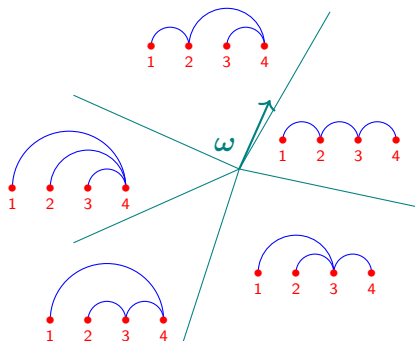
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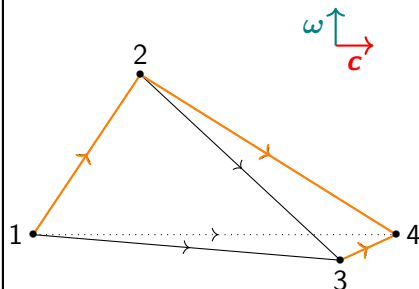
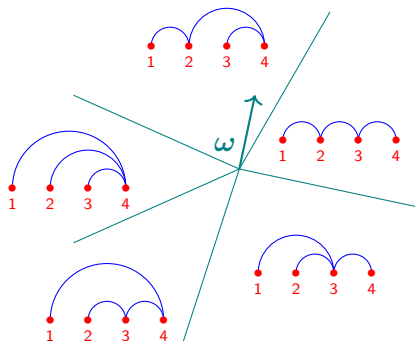
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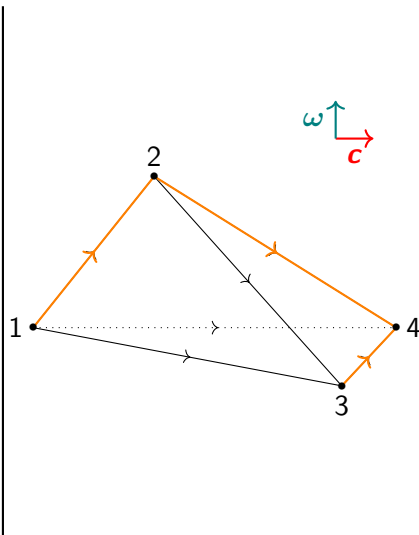
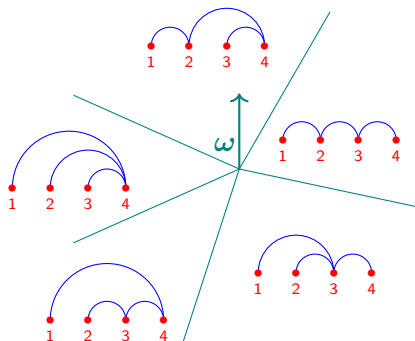
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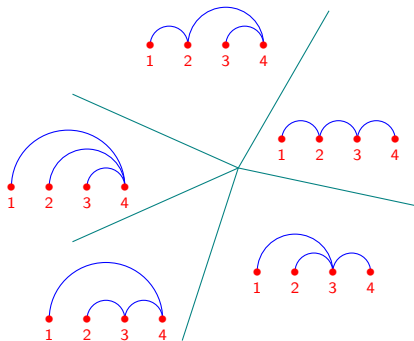
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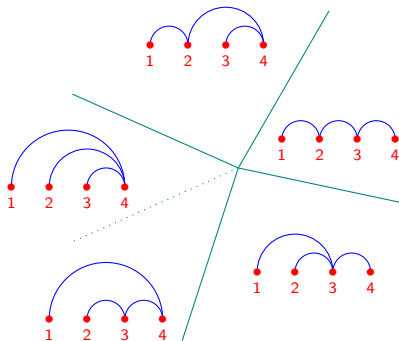
(Max-slope) pivot polytope



Coherent arborescence:
arborescence obtained via
max-slope pivot rule

Pivot rule fan:
 $\omega \sim \omega'$ iff same arborescence.

(Max-slope) pivot polytope



Coherent arborescence:

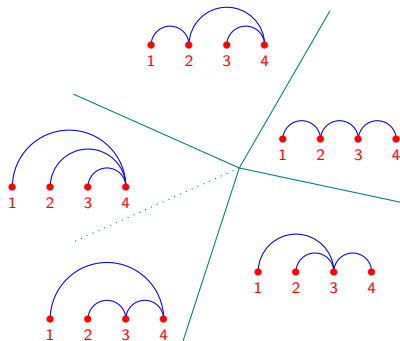
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Pivot rule fan refines monotone
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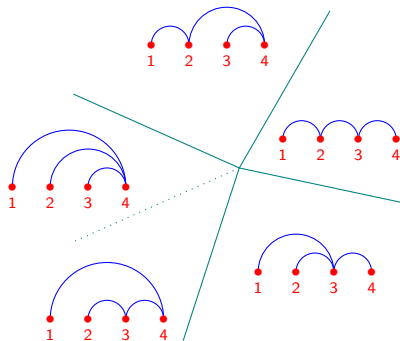
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Theorem (Black, De Loera, Lütjeharms, Sanyal '22)

The pivot rule fan is polytopal.

(Max-slope) pivot polytope Π_c : dual to the pivot rule fan

(Max-slope) pivot polytope



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$$\Sigma_c(\Delta_d) = \text{Cube}_{d-1}$$

$$\Pi_c(\Delta_d) = \text{Asso}_d$$

Link with generalized permutahedra?

Theorem (Black, De Loera, Lütjeharms, Sanyal '23+)

$$\Pi_c(\Delta_d) \simeq \text{Asso}_d \quad \text{for all } \mathbf{c}$$

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Conjecture (Pilaud, Sanyal)

$$\Pi_c(\Delta_{d_1} \times \Delta_{d_2}) \simeq \text{Asso}_{d_1} \star \text{Asso}_{d_2} \quad \star \text{ shuffle product}$$

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Challenge 1: Prove the conjecture!

Link with generalized permutahedra?

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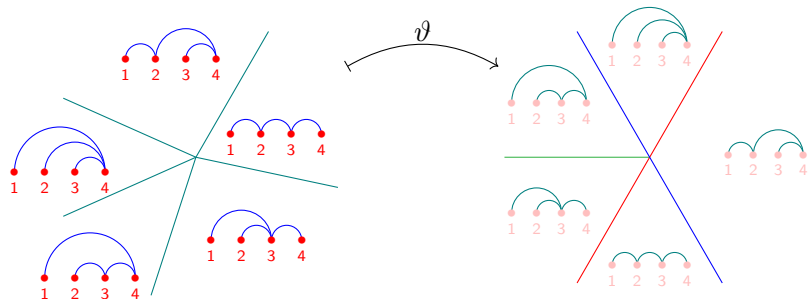
Conjecture (Pilaud, Sanyal)

$$\Pi_c(\Delta_{d_1} \times \Delta_{d_2}) \simeq \text{Asso}_{d_1} \star \text{Asso}_{d_2} \quad \star \text{ shuffle product}$$

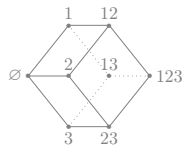
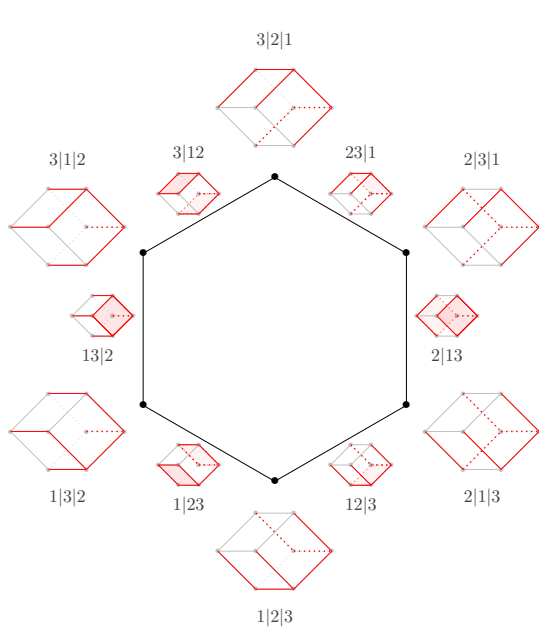
Challenge 1: Prove the conjecture!

Challenge 2: Give **geometric** proofs!

Case of the d -simplex



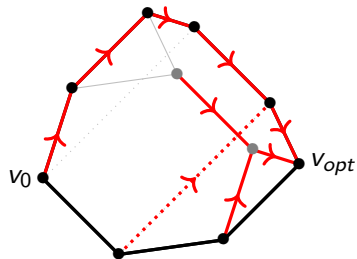
Case of the d -cube



Slope comparisons

Max-slope pivot rule: take (improving)
neighbor with best slope

For ω , what is important?

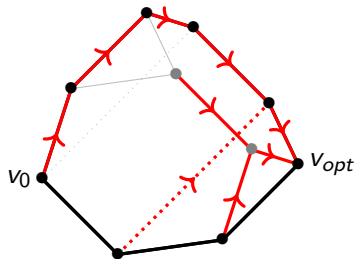


Slope comparisons

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Slopes: $\tau_{\omega}(u, v) = \frac{\langle \omega, u-v \rangle}{\langle c, u-v \rangle}$



Slope comparisons

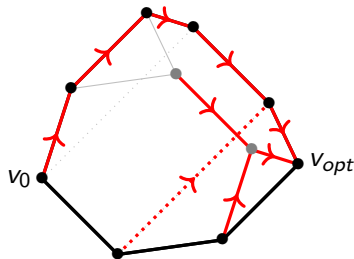
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Slope vector:

$$\theta(\omega) = (\tau_\omega(u, v))_{uv \in E(P)}$$



Slope comparisons

Max-slope pivot rule: take (improving)
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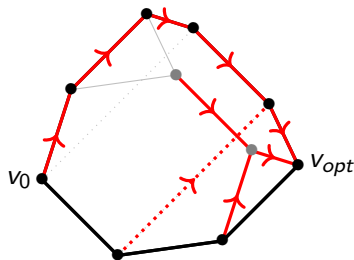
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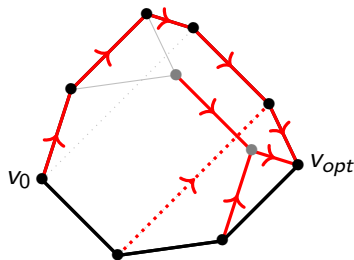
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What is **really** important??



Slope comparisons

Max-slope pivot rule: take (improving)
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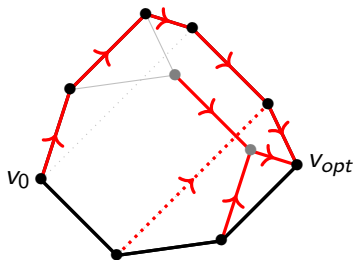
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What is **really** important?? The comparisons of slopes!



Slope comparisons

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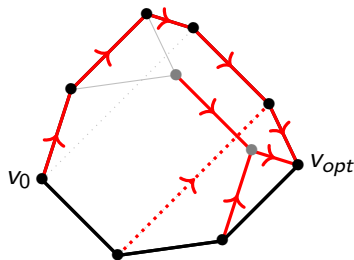
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What is **really** important?? The comparisons of slopes!

Compare coordinates of $\theta(\omega)$



Slope comparisons

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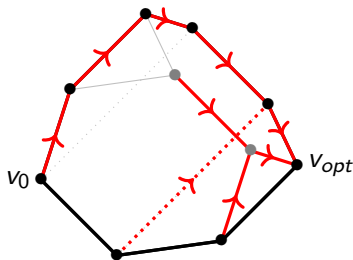
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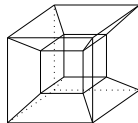
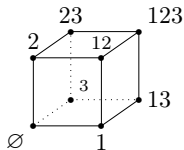
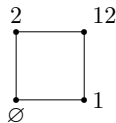
What is **really** important?? The comparisons of slopes!

Compare coordinates of $\theta(\omega)$

Where is $\theta(\omega)$ in the braid fan $\mathcal{B}_m$?

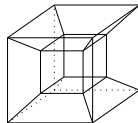
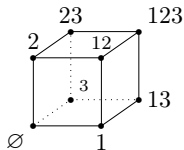
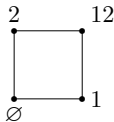


Case of the d -cube



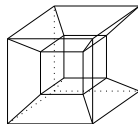
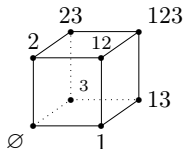
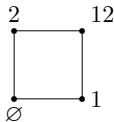
Too many edges

Case of the d -cube



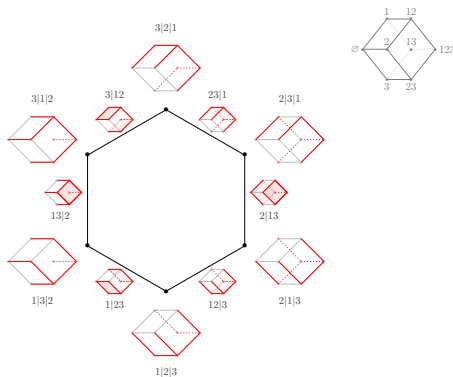
Too many edges, **but**
parallelism saves us!

Case of the d -cube

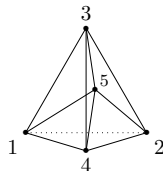
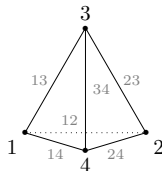
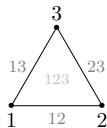
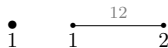


Too many edges, **but**
parallelism saves us!

Geometric proof of
 $\Pi_c(\text{Cube}_d) = \Pi_d$

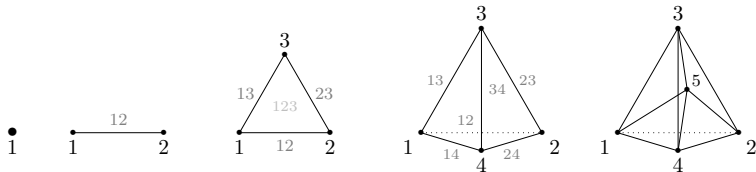


Case of the d -simplex



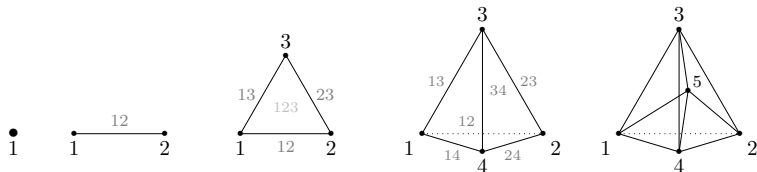
Too many edges

Case of the d -simplex



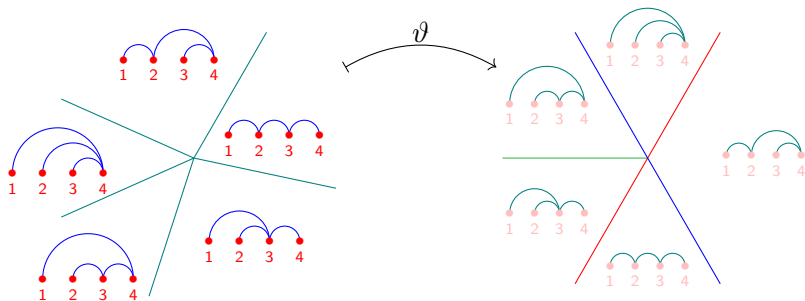
Too many edges, **but** affine independence saves us!

Case of the d -simplex



Too many edges, **but** affine independence saves us!

Geometric proof of $\Pi_c(\Delta_d) = \text{Asso}_d$



Shuffle: $(E, \leq)$ and $(F, \preceq)$ posets, then $\trianglelefteq$ is a shuffle when:

ground set : $E \sqcup F$

relations : all relations of $\leq$; all relations of $\preceq$;

for each $e \in E, f \in F$, choose if $e \trianglelefteq f$ or $e \triangleright f$
(+ transitive closure)

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(+ transitive closure)

Theorem (Chapoton, Pilaud '22)

P, Q : *generalized permutahedra*.

Exists polytope $P \star Q$ s.t.

$$\mathcal{P}(P \star Q) = \{ \text{all shuffles between } \leq \in \mathcal{P}(P) \text{ and } \preceq \in \mathcal{P}(Q) \}$$

Combine parallelism & affine independence:

Theorem

For $\Delta_{d_1} \times \cdots \times \Delta_{d_r}$, all (generic) direction:

$$\Pi_c(\Delta_{d_1} \times \cdots \times \Delta_{d_r}) \simeq \text{Asso}_{d_1} \star \cdots \star \text{Asso}_{d_r}$$

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Example

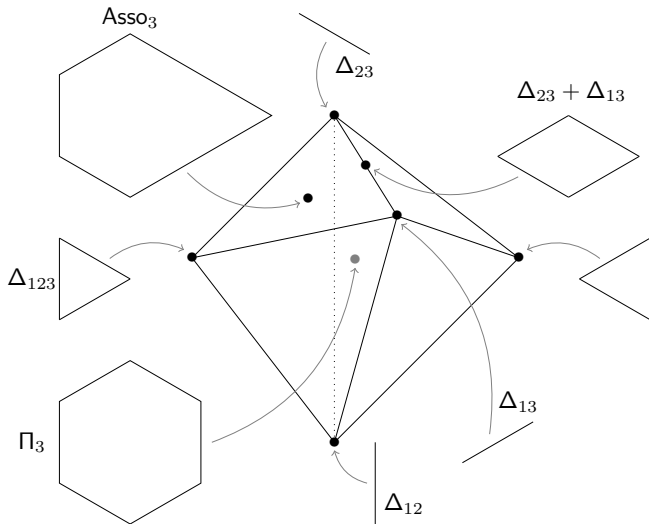
- (a) $\Pi_c(\square_d) \simeq \text{Perm}_d$
- (b) $\Pi_c(\square_m \times \Delta_n) \simeq (m, n)\text{-multiplihedron}$
- (c) $\Pi_c(\Delta_m \times \Delta_n) \simeq (m, n)\text{-constrainahedron}$

- 1) For which P , $\Pi_c(P)$ is a generalized permutahedron?
→ a priori, only products of simplices, but no proof
- 2) Is $\Pi_c(P)$ projection of a generalized permutahedron?
→ pivot fan sent inside $\text{Im}(\theta) \cap \mathcal{B}_m$
- 3) When $\Pi_c(P)$ and $\Pi_c(Q)$ **not** generalized permutahedra, what happen to $\Pi_c(P \times Q)$?
→ not equivalent to $\Pi_c(P) \star \Pi_c(Q)$, but "embeds" in it

Contents

Introduction	5
1 Preliminaries	11
1.1 Partially ordered sets	11
1.2 Polytopes	12
1.2.1 Simplex	15
1.2.2 Cube	15
1.2.3 Permutohedron	15
1.2.4 Associahedron	18
1.3 Linear programming	20
2 Deformations of polytopes and generalized permutohedra	24
2.1 Deformations of polytopes	24
2.2 Deformation cones of graphical zonotopes	28
2.2.1 Graphical zonotopes	28
2.2.2 Graphical deformation cones	29
2.2.3 The facets of graphical deformation cones	33
2.2.4 Simplicial graphical deformation cones	36
2.2.5 Perspectives and open questions	37
2.3 Deformation cones of nestohedra	38
2.3.1 Deformation cones of graphical nested fans	39
2.3.2 Deformation cones of arbitrary nested fans	47
2.3.3 Simplicial deformation cones and interval building sets	61
2.3.4 Perspectives and open questions	63
3 Max-slope pivot rule polytopes	65
3.1 Max-slope pivot rule and max-slope pivot polytope	65
3.2 Max-slope pivot polytope of cyclic polytopes	70
3.2.1 Cyclic associahedra and the intrinsic degree	71
3.2.2 Realization sets and universal arborescence	74
3.2.3 Pivot polytopes of cyclic polytopes of dimension 2 and 3	86
3.2.4 Perspectives and open questions	92
3.3 Max-slope pivot polytopes of products of polytopes	95
3.3.1 Max-slope pivot polytopes of the cube and the simplex	99
3.3.2 Max-slope pivot polytope of a product of simplices	102
3.3.3 Perspectives and open questions	106
4 Fiber polytopes	108
4.1 Preliminaries on fiber polytopes	108
4.2 Monotone path polytopes of the hypersimplices	111
4.2.1 Monotone paths polytopes in general	111
4.2.2 A necessary criterion for coherent paths on $\Delta(n, k)$	119
4.2.3 Sufficiency of this criterion in the case $\Delta(n, 2)$	119
4.2.4 Counting the number of coherent monotone paths on $\Delta(n, 2)$	127
4.2.5 Perspectives and open questions	130
4.3 Fiber polytopes for the projection from $\text{Cyc}_d(t)$ to $\text{Cyc}_2(t)$	133
4.3.1 Bijection between triangulations and non-crossing arborescences	133
4.3.2 Fiber polytopes for the projection $\text{Cyc}_d(t) \xrightarrow{\pi} \text{Cyc}_2(t)$	135
4.3.3 Realization sets and universal triangulations for $\text{Cyc}_d(t) \xrightarrow{\pi} \text{Cyc}_2(t)$	138
4.3.4 Perspectives and open questions	143
A A Vandermonde-like determinant	146

Thank you!



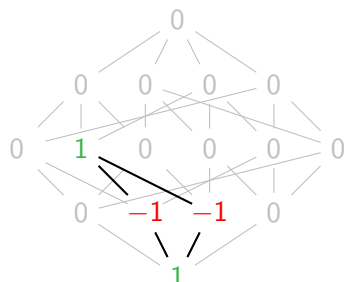
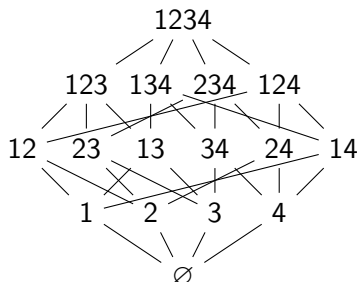
Bonus slides...

The tool: submodular dependancies

Notations: $Sx = S \cup \{x\}$, $(\mathbf{f}_x)_{x \subseteq [n]}$ canonical basis of $\mathbb{R}^{2^{[n]}}$

Definition

Submodular vector $\mathbf{n}(S, u, v) = \mathbf{f}_{Suv} - \mathbf{f}_{Su} - \mathbf{f}_{Sv} + \mathbf{f}_S$
for $u, v \in S \subseteq [n]$

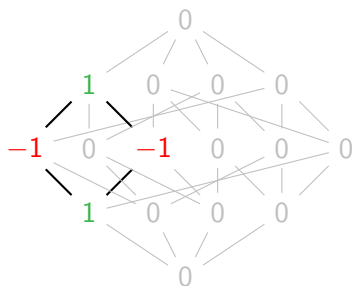
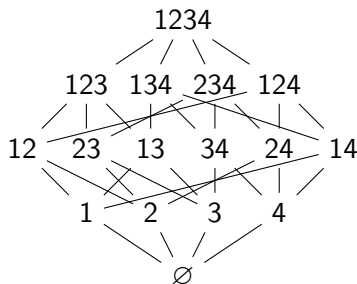


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The tool: submodular dependancies

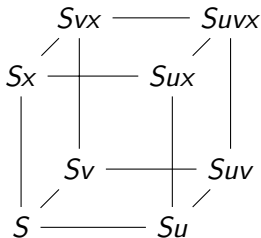
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$\mathbf{n}(S, u, v)$ are the facet's normals of $\mathbb{DC}(\Pi_n)$

Lemma (Cubic relation)

$u, v, x \notin S \subseteq [n]$

$$\mathbf{n}(Suvx, u, v) + \mathbf{n}(Sux, u, x) = \mathbf{n}(Suv, u, v) + \mathbf{n}(Suvx, u, x)$$



NB: Cubic relations generates all relations of submodular vectors

The tool: submodular dependancies

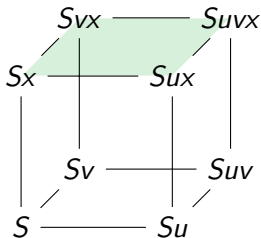
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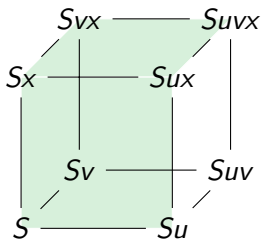
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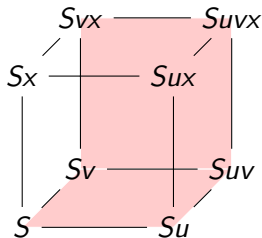
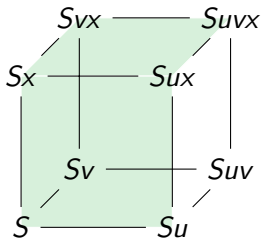
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NB: Cubic relations generates all relations of submodular vectors

Monotone path polytope and pivot rule polytope

Let $P \subset \mathbb{R}^d$ be a polytope.

Max-slope pivot rule: $A^\omega(v) = \operatorname{argmax} \left\{ \frac{\langle \omega, u-v \rangle}{\langle c, u-v \rangle}; u \text{ impr. neg. of } v \right\}$.

Coherent monotone path: A monotone path that can be obtained via the max-slope pivot rule.

Monotone path polytope $\Sigma_c(P)$ [?]: Fiber polytope of $P \xrightarrow{\pi} Q$ with Q a segment. (Can be seen as a Minkowski sum of sections of P .)

The vertices of $\Sigma_c(P)$ are all coherent monotone paths.

Coherent arborescence: An arborescence that can be obtained via the max-slope pivot rule.

Pivot rule polytope $\Pi_c(P)$: Polytope which vertices are all coherent arborescences.

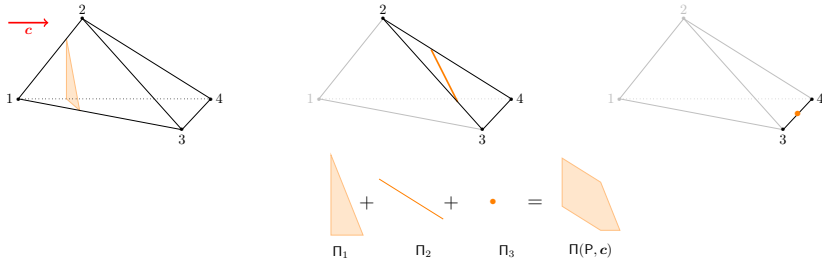
$$\Pi_c(P) = \operatorname{conv} \left\{ \sum_{v \neq v_{\text{opt}}} \frac{1}{\langle c, A(v) - v \rangle} (A(v) - v); A \text{ coherent arbo. of } P \right\}$$

Monotone path polytope and pivot rule polytope

Coherent arborescence: An arborescence that can be obtained via the max-slope pivot rule.

Pivot rule polytope $\Pi_c(P)$: Polytope whose vertices are all coherent arborescences. Can also be seen as a Minkowski sum of sections:

$$\sum_{v \in V(P)} (\text{section between } v \text{ and its improving neighbors})$$



Mimicking the case of the d -cube

Idea 1:

Fix a polytope P , and direction $\mathbf{c}$, n vertices, m edges.

$\theta : \mathbb{R}^d \rightarrow \mathbb{R}^m$ sends the pivot fan inside $\text{Im}(\theta) \cap \mathcal{B}_m$

Problem: **This** is not a braid fan as $d \ll m \dots$

If m' classes of parallelism:

$\bar{\theta} : \mathbb{R}^d \rightarrow \mathbb{R}^{m'}$ sends the pivot fan inside $\text{Im}(\bar{\theta}) \cap \mathcal{B}_{m'}$

Problem: **This** is not a braid fan as $d \ll m' < m \dots$

We need to go lower dimensional!

Adapted slope map

Idea 2:

Fix a polytope P , direction $\mathbf{c}$,
 n vertices, m edges.

Fix A arborescence:

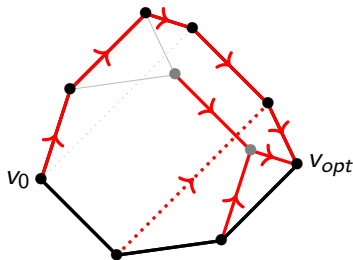
$$\vartheta_A(\omega) = (\tau_\omega(u, A(u)) ; u \text{ vertex})$$

ϑ_A : linear, injective, $\mathbb{R}^d \rightarrow \mathbb{R}^{n-1}$

but if ω does not capture A , then $\vartheta_A(\omega)$ have no meaning...

Adapted slope map: $\vartheta(\omega) = \vartheta_{A^\omega}(\omega)$

i.e. take ω and look at the slope of the edges it selects.



Case of the d -simplex

$d = n - 1 \iff P$ is a simplex

For Δ_d : $\vartheta : \mathbb{R}^d \rightarrow \mathbb{R}^d$ piece-wise linear, $\ker \vartheta = \{\mathbf{0}\} \Rightarrow$ bijection
 ϑ sends the pivot fan of Δ_d inside $\mathcal{B}_d$.

